

## A Appendix: Proofs

### A.1 Selected Observations

ad h):  $a \in \mathbf{RT}$  iff  $a^* = a$ .

We reason by induction on  $|a|$ , where  $a \in \mathbf{FT}$ . If  $|a| = 1$ , then  $a \in \{t, e\}$  and hence  $a \in \mathbf{RT}$  and  $a^* = a$  by definition. For the inductive step, we separate the directions. “ $\Rightarrow$ ”: There are two cases: either (a)  $a = (st)$ ; or (b)  $a = (bc)$ , where  $b, c \in \mathbf{RT}$ . If (a),  $a^* = a$ , again by definition; if (b),  $b = b^*$  and  $c = c^*$ , by induction. Moreover,  $(bc)^* = (b^*c^*)$ , by the definition of the shift  $*$  and given that  $b = b^*$ . Hence  $a^* = (b^*c^*) = (bc)$ . – “ $\Leftarrow$ ”: If  $|a| > 1$  and  $a \in \mathbf{FT}$ , then there are two cases: either (c)  $a = (bc)$ , or (d)  $a = (sc)$ , where  $b, c \in \mathbf{FT}$ ; in both cases, we reason by contraposition. Turning to (c), an inspection of the definition of  $*$  reveals that  $a^* = (b^*d)$ , for some  $d \in \mathbf{2T}$ . However, if  $a \notin \mathbf{RT}$ , then either (c1)  $b \notin \mathbf{RT}$  or else (c2)  $c \notin \mathbf{RT}$ , given the definition of  $\mathbf{RT}$ . But in case (c1),  $b \neq b^*$ , by induction, and thus  $a^* = (b^*d) \neq (bc) = a$ , by the characteristic property of ordered pairs.<sup>i</sup> Yet if (c2) applies,  $c \neq c^*$ , by induction, and at the same time,  $d = c^*$  or  $d = (c^*t)$ , depending on whether  $c^*$  is Boolean. Either way,  $d \neq c : c \neq c^*$  by assumption, and  $c \neq (c^*t)$  because  $|c^*t| > |c|$ , by **f**). So,  $a^* = (b^*d) \neq (bc) = a$ . – As to case (d), the two basic sub-cases are immediate:  $(se)^* = (e(st)) \neq (se)$ ; and  $(st) \in \mathbf{RT}$ , anyway. In the remaining two cases,  $a^* = (bd)$ , for certain  $b, d \in \mathbf{FT}$  – and thus:  $a^* = (bd) \neq (sc)$ , because  $s \notin \mathbf{FT}$ .

ad j): If  $a \notin \mathbf{RT}$ ,  $a^* \in \mathbf{BT}$ .

An induction on  $a \in \mathbf{FT}$  reveals that:

**j\*)** If  $a^* \notin \mathbf{BT}$ ,  $a^* = a$ ,

which is equivalent to **j**), in view of **h**). More specifically, if  $a \in \{e, t\}$ , **j\*)** holds trivially. If  $a = (bc)$  and  $a^* \notin \mathbf{BT}$ , then  $a^* = (b^*c^*)$ , where  $c^* \notin \mathbf{BT}$ . But then  $b = b^*$ , by the definition of  $(bc)^*$ ; moreover, by induction,  $c^* = c$ . So  $a^* = (b^*c^*) = (bc) = a$ . Finally, if  $a = (sb)$ ,  $a^* \in \mathbf{BT}$ , by **i**).

ad m):  $\dot{\kappa}^\cup$  is (i) surjective but (ii) not injective.

It should be noted that  $\dot{\kappa}^\cup: \bigcup_{a \in \mathbf{FT}} D_a \longrightarrow \bigcup_{a \in \mathbf{RT}} D_a$  only because we have been assuming that  $(+)$   $D_a \cap D_b = \emptyset$  whenever  $a \neq b$ ; otherwise  $\dot{\kappa}^\cup$  might not even be a function to begin with. But given  $(+)$ , (i) immediately follows from **l**): if  $a \in \mathbf{RT}$  and  $y \in D_b$ , then  $y = id_b(y) = \dot{\kappa}_b(y) = \dot{\kappa}^\cup(y)$ . As to (ii), whenever  $a \in \mathbf{FT} \setminus \mathbf{RT}$  and  $x \in D_a$ ,  $\dot{\kappa}_a(x) = y \in D_{a^*}$ , where  $a^* \in \mathbf{RT}$ . Hence, by **l**),  $\dot{\kappa}_{a^*}(y) = id_{a^*}(y) = y$ . But then:  $\dot{\kappa}^\cup(x) = \dot{\kappa}_a(x) = y = \dot{\kappa}_{a^*}(y) = \dot{\kappa}^\cup(y)$ , even though  $x \neq y$ , due to the disjointness assumption  $(+)$ .

<sup>i</sup>However ordered pairs are defined, they satisfy the condition that  $(x, y) = (x', y')$  iff  $x = x'$  and  $y = y'$ .

ad n):  $g[x/u]^* = g^*[\tilde{x}/\dot{\kappa}_a(u)]$

To begin with, we have:

$$\begin{aligned} g[x/u]^*(\tilde{x}) &= \\ \dot{\kappa}_a(g[x/u](x)) &= \text{by the definition of } h^* \\ \dot{\kappa}_a(u) &= \text{by the definition of } h[y/v] \\ g^*[\tilde{x}/\dot{\kappa}_a(u)](\tilde{x}) &= \text{by the same definition} \end{aligned}$$

But then, for any  $y = z^* \in Var \setminus \{\tilde{x}, i\}$ :

$$\begin{aligned} g[x/u]^*(y) &= \\ \dot{\kappa}_a(g[x/u](z)) &= \text{by the definition of } h^* \\ \dot{\kappa}_a(g(z)) &= y \neq \tilde{x} \text{ and thus } z \neq x \\ g^*(y) &= \text{by the definition of } h^* \\ g^*[\tilde{x}/\dot{\kappa}_a(u)](y) &= y \neq \tilde{x} \end{aligned}$$

ad q): The two claims being very similar, we only prove the second one:

$$\overline{\sqcap}_b(/X/) = \mathbf{X}.$$

Given  $b = (\vec{b} \text{ t}) \in \mathbf{RT} \cap \mathbf{BT}$ , where  $\vec{b} = \langle b_1, \dots, b, s \rangle$ ,  $\mathbf{X} \in D_b$  and arbitrary  $\mathbf{x}_1 \in D_{b_1}, \dots, \mathbf{x}_n \in D_{b_n}$  and  $\mathbf{w} \in D_s$ , we have:

$$\begin{aligned} &\overline{\sqcap}_b(/X/)(\mathbf{x}_1) \dots (\mathbf{x}_n)(\mathbf{w}) = 1 \\ \text{iff } &\mathbf{P}(\mathbf{x}_1) \dots (\mathbf{x}_n)(\mathbf{w}) = 1, \text{ for any } \mathbf{P} \in D_b \text{ such that } /X/(\mathbf{P}) = 1 \quad \text{IL}^*\text{-int.} \\ \text{iff } &\mathbf{P}(\mathbf{x}_1) \dots (\mathbf{x}_n)(\mathbf{w}) = 1, \text{ for any } \mathbf{P} \in D_b \text{ such that } \mathbf{P} = \mathbf{X} \quad \text{def. } / \cdot / \\ \text{iff } &\mathbf{X}(\mathbf{x}_1) \dots (\mathbf{x}_n)(\mathbf{w}) = 1, \quad \text{identity} \end{aligned}$$

which means that  $\overline{\sqcap}_b(/X/) = \mathbf{X}$ , by extensionality.

ad t):  $\dot{\kappa}_a$  is  $(aa^*)$ -invariant .

Let  $(\pi_a)_{a \in \mathbf{2T}}$  be a permutation of  $(D_a)_{a \in \mathbf{2T}}$  (where  $D_e$  and  $D_s$  are arbitrary, non-empty sets). Then for any  $IL^*$ -model  $F$  and  $IL^*$ -assignment  $g$ , the corresponding *permutation model*  $F_\pi$  and *permutation assignment*  $g_\pi$  are defined as:  $F_\pi := \pi \circ F$  and  $g_\pi := \pi \circ g$ , where ‘ $\circ$ ’ denotes functional composition.<sup>ii</sup> A routine induction shows that, for any  $a \in \mathbf{FT}$  and any  $\alpha \in IL_a^* : (*)$   $\llbracket \alpha \rrbracket^{F_\pi, g_\pi} = \pi(\llbracket \alpha \rrbracket^{F, g})$ . But then:

$$\begin{aligned} &\dot{\kappa}_a \\ = &\llbracket \kappa_a \rrbracket^{F, g} \quad \text{def. } \dot{\kappa}_a \\ = &\llbracket \kappa_a \rrbracket^{F_\pi, g_\pi} \quad \text{Coincidence Lemma}^{\text{iii}} \\ = &\pi(\llbracket \kappa_a \rrbracket^{F, g}) \quad (*) \\ = &\pi(\dot{\kappa}_a) \quad \text{def. } \dot{\kappa}_a \end{aligned}$$

The logicity of  $aa^*$  and the embedability of  $a$  into  $a^*$  then immediately follow from the  $\kappa$ -Lemma and the pertinent definitions.

<sup>ii</sup>In other words,  $F_\pi(c) = \pi(F(c))$  and  $g_\pi(x) = \pi(g(x))$  for any  $c \in Con$  and  $x \in Var$ .

<sup>iii</sup>The very  $\dot{\kappa}$ -notation had already taken this step for granted. Strictly speaking, it requires a more general (and standard) version of the Coincidence Lemma that generalizes over models as well as sets of individuals and words.

ad  $\mathbf{u}$ :  $((se)e)((se)e)^o = ((se)e)((e(st))e)$  is not logical.

To be sure,  $((se)e)^o = (se)^oe^o = (e^o(st))e^o = ((e(st))e)$ , by the very definition of Kaplan's shift  $^o$ . Taking advantage of the model-theoretic perspective, we assume that  $D_s = \{\mathbf{w}_0\}$  is a singleton, whereas  $D_e$  has at least two members, and show that for any  $\mathbf{F} \in D_{((se)e)((e(st))e)}$  there is a permutation  $\pi$  of  $(D_a)_{a \in 2\mathbf{T}}$  such that  $\pi_{((se)e)((e(st))e)}(\mathbf{F}) \neq \mathbf{F}$ . To adopt a strategem of van Benthem [1, p. 319], we first observe that, given the circumstances, the domains  $D_e$  and  $D_{se}$  are isomorphic in that the elements of  $D_{se}$  are of the form  $\mathbf{c}_x := \{(\mathbf{w}_0, \mathbf{x})\}$ , where  $\mathbf{x} \in D_e$  and  $\mathbf{c}_x \neq \mathbf{c}_y$  for distinct  $\mathbf{x}, \mathbf{y} \in D_e$ . Moreover, for any permutation  $\pi$  and any  $\mathbf{c}_x \in D_{se}$  we have:

$$\begin{aligned}
(*) \quad & \pi_{se}(\mathbf{c}_x)(\mathbf{w}_0) \\
= & \pi_e(\mathbf{c}_x(\pi_s^{-1}(\mathbf{w}_0))) && \text{def. } \pi_{se} \\
= & \pi_e(\mathbf{c}_x(\mathbf{w}_0)) && D_s = \{\mathbf{w}_0\} \\
= & \pi_e(\mathbf{x}) && \text{def. } \mathbf{c}_x \\
= & \mathbf{c}_{\pi_e(\mathbf{x})}(\mathbf{w}_0) && \text{def. } \mathbf{c}_{\pi_e(\mathbf{x})}
\end{aligned}$$

In other words:

$$(*) \quad \pi_{se}(\mathbf{c}_x) = \mathbf{c}_{\pi_e(\mathbf{x})},$$

by extensionality and given that  $D_s = \{\mathbf{w}_0\}$ . As a consequence, the function  $\mathbf{g} := \{(\mathbf{c}_x, \mathbf{x}) \mid \mathbf{x} \in D_e\}$  is  $((se)e)$ -invariant; for if  $\pi$  is a permutation and  $\mathbf{c}_x \in D_{se}$ , we have:

$$\begin{aligned}
& \pi_{(se)e}(\mathbf{g})(\mathbf{c}_x) \\
= & \pi_e(\mathbf{g}(\pi_{se}^{-1}(\mathbf{c}_x))) && \text{def. } \pi_{(se)e} \\
= & \pi_e(\mathbf{g}(\mathbf{c}_{\pi_e^{-1}(\mathbf{x})})) && \text{by } (*) \\
= & \pi_e(\pi_e^{-1}(\mathbf{x})) && \text{def. } \mathbf{g} \\
= & \mathbf{x} && \pi_e \circ \pi_e^{-1} = id_e \\
= & \mathbf{g}(\mathbf{c}_x) && \text{def. } \mathbf{g}
\end{aligned}$$

Now suppose  $\mathbf{F} \in D_{((se)e)((e(st))e)}$  and put:  $\mathbf{x}_0 := \mathbf{F}(\mathbf{g})(\mathbf{0}_{e(st)}) \neq \mathbf{x}_1 \in D_e$ . We may then consider the transposition

$$\pi^\sharp := [id_e \setminus \{(\mathbf{x}_0, \mathbf{x}_0), (\mathbf{x}_1, \mathbf{x}_1)\}] \cup \{(\mathbf{x}_0, \mathbf{x}_1), (\mathbf{x}_1, \mathbf{x}_0)\}$$

and observe:

$$\begin{aligned}
& \pi_{((se)e)((e(st))e)}^\sharp(\mathbf{F})(\mathbf{g})(\mathbf{0}_{e(st)}) \\
= & \pi_{(e(st))e}^\sharp(\mathbf{F}([\pi_{(se)e}^\sharp]^{-1}(\mathbf{g}))) (\mathbf{0}_{e(st)}) && \text{def. } \pi_{((se)e)((e(st))e)}^\sharp \\
= & \pi_{(e(st))e}^\sharp(\mathbf{F}(\mathbf{g})) (\mathbf{0}_{e(st)}) && \mathbf{g}'\text{'s invariance} \\
= & \pi_e^\sharp(\mathbf{F}(\mathbf{g})([\pi_{e(st)}^\sharp]^{-1}(\mathbf{0}_{e(st)}))) && \text{def. } \pi_{(e(st))e}^\sharp \\
= & \pi_e^\sharp(\mathbf{F}(\mathbf{g})(\mathbf{0}_{e(st)})) && \mathbf{0}_{e(st)}\text{'s invariance} \\
= & \pi_e^\sharp(\mathbf{x}_0) && \text{def. } \mathbf{x}_0 \\
= & \mathbf{x}_1 && \text{def. } \pi_e^\sharp \\
\neq & \mathbf{x}_0 && \text{assumption} \\
= & \mathbf{F}(\mathbf{g})(\mathbf{0}_{e(st)}) && \text{def. } \mathbf{x}_0
\end{aligned}$$

Hence  $\pi_{((se)e)((e(st))e)}^\sharp(\mathbf{F}) \neq \mathbf{F}$ , which means that  $\mathbf{F}$  is not invariant.

ad w): (I)  $a = a^*$  iff  $a = a^+$ , and (II)  $a^* \in \mathbf{BT}$  iff  $a^+ \in \mathbf{BT}$ .

We reason by induction on  $a \in \mathbf{FT}$ . If  $a = t$  or  $a = e$ ,  $a^* = a = a^+$  and thus both (I) and (II) hold trivially. – If  $a = (bc)$  and  $a^* = (b^*c^*)$ , then either:  $b = b^*$ , or:  $c^* \in \mathbf{BT}$ , and so either:  $b = b^+$ , or  $b^+ \in \mathbf{BT}$ , by induction. Hence  $a^+ = (b^+c^+)$  and we have:

|      |                         |                                          |
|------|-------------------------|------------------------------------------|
| (I)  | $a = a^*$               |                                          |
| iff  | $b = b^*$ and $c = c^*$ | characteristic property of ordered pairs |
| iff  | $b = b^+$ and $c = c^+$ | induction                                |
| iff  | $a = a^+$               | characteristic property of ordered pairs |
| (II) | $a^* \in \mathbf{BT}$   |                                          |
| iff  | $c^* \in \mathbf{BT}$   | def. $\mathbf{BT}$                       |
| iff  | $c^+ \in \mathbf{BT}$   | induction                                |
| iff  | $a^+ \in \mathbf{BT}$   | def. $\mathbf{BT}$                       |

If  $a = (bc)$  and  $a^* \neq (b^*c^*)$ , then  $a^* = (b^*(c^*t))$ ,  $b \neq b^*$ , and  $c^* \notin \mathbf{BT}$ , by the definition of  $(bc)^*$ . Hence, by induction:  $b \neq b^+$ , and  $c^+ \notin \mathbf{BT}$ , whence  $a^+ = (b^+c^+(t))$ . So (I) holds because  $a = (bc) \neq a^* = (b^*(c^*t))$ , given that  $|c| \leq |c^*| < |c^*| + 1 = |(c^*t)|$ , and  $a = (bc) \neq a^+ = (b^+(c^+t))$ , by the same token.<sup>iv</sup> And (II) holds given that both  $a^*$  and  $a^+$  are Boolean. – If  $a = (st)$ ,  $a^* = a = a^+$  and thus both (I) and (II) hold trivially. – If  $a = (se)$ ,  $a^* = (e(st)) = (e^+(st)) = a^+$  and again both (I) and (II) hold trivially. – If  $a = (s(bc))$ , where  $b \neq s$ , then  $a^* = (b^*(sc^*)) \neq (s(bc)) = a$ ; for otherwise we would have:  $s = b^* \in \mathbf{BT}$ . But then  $a^+ = ((bc)^+(st)) \neq (s(bc)) = a$ ; for otherwise  $s = (bc)^+ \in \mathbf{RT}$ . Hence  $a$  satisfies (I). Moreover,  $a^* \in \mathbf{BT}$ , by (i), and  $a^+ = ((bc)^+(st)) \in \mathbf{BT}$ , which settles (II). – Finally (and similarly), if  $a = (s(sb))$ , then  $a^* = (sb^*(sb^*)(st)) \neq a$ ; for otherwise we would have:  $s = (sb)^* \in \mathbf{BT}$ . But then  $a^+ = ((sb)^+(st)) \neq (s(sb)) = a$ ; for otherwise  $s = (sb)^+ \in \mathbf{RT}$ . Hence  $a$  satisfies (I). Moreover,  $a^* \in \mathbf{BT}$ , by (i), and  $a^+ = ((sb)^+(st)) \in \mathbf{BT}$ , which settles (II) and completes the proof.

## A.2 Footnotes

### A.2.1 Fn. 14

We need to find a bijection  $x \mapsto \tilde{x}$  between  $\bigcup_{a \in \mathbf{FT}} \text{Var}_{a \setminus \{i\}}$  and  $\bigcup_{a \in \mathbf{RT}} \text{Var}_a$  such that, for any  $a \in \mathbf{FT}$ ,  $\tilde{x} \in \text{Var}_{a^*}$  if  $x \in \text{Var}_a$ . To construct such  $\tilde{\cdot}$ , we represent all variables of a type  $a \in \mathbf{FT}$  as  $v_{n,a}$ , where  $n \in \mathbb{N}$ . First, we observe that, for any  $c \in \mathbf{RT}$ , the set  $[c] := \{a \in \mathbf{FT} \mid a^* = c\}$  is finite, by **f**), and non-empty, by **h**), and can thus be represented by way of a (repetition-free) list of some length  $|c| + 1$ :  $[c] = \{[c]_0, \dots, [c]_{|c|}\} \neq \emptyset$ . As a case in point,  $[e(st)] = \{e(st), s(et), se\}$ , whence  $|e(st)| = 2$ . Keeping the enumeration fixed

<sup>iv</sup>We leave it to the reader to show that, for any  $a \in \mathbf{BT}$ ,  $|a| \leq |a^+|$ , which is just as straightforward as observation **f**).

(for each  $c \in \mathbf{RT}$ ), any type  $a \in \mathbf{FT}$  is thus identifiable by its position  $\#a$ :  $a = [a^*]_{\#a}$ , where,  $\#a \neq \#b$  if  $a^* = b^*$  and  $a \neq b$ . Continuing the example, we thus have:  $\#se = 2$ . We now put, for any  $n \in \mathbb{N}$  and  $a \in \mathbf{FT}$ :

$$\widetilde{v_{n,a}} := v_{f(n,a),a^*}, \text{ where } f(n,a) = n \times (\backslash a^* \backslash + 1) + \#a.$$

Using the same enumeration as before, we thus have:  $\widetilde{v_{2,se}} = v_{8,e(st)}$ ; and the variables of type  $a = e(st)$  are arranged as follows, where  $b = s(et)$ , and  $c = se$ :  $v_{0,a} = \widetilde{v_{0,a}}$ ;  $v_{1,a} = \widetilde{v_{0,b}}$ ;  $v_{2,a} = \widetilde{v_{0,c}}$ ;  $v_{3,a} = \widetilde{v_{1,a}}$ ;  $v_{4,a} = \widetilde{v_{1,b}}$ ;  $v_{5,a} = \widetilde{v_{1,c}}$ ; etc.

Obviously,  $\tilde{x} \in Var_{a^*}$  whenever  $x \in Var_a$  and  $a \in \mathbf{FT}$ . It remains to be shown that  $x \mapsto \tilde{x}$  is (a) one-one and (b) onto  $\bigcup_{a \in \mathbf{RT}} Var_a$ . ad (a): if  $x, y \in Var \setminus \{i\}$  and  $x = v_{n,a} \neq y = v_{m,b}$ , where  $a^* \neq b^*$ ,  $\tilde{x} \in Var_{a^*}$  and  $\tilde{y} \in Var_{b^*}$ , and so  $\tilde{x} \neq \tilde{y}$ , since  $Var_{a^*} \cap Var_{b^*} = \emptyset$ .<sup>v</sup> But if  $a^* = b^*$ , and  $x = v_{n,a} \neq v_{m,b} = y$ , then either (a1)  $\#a = \#b$ , or (a2)  $\#a \neq \#b$  but  $m \neq n$ . In either case we have  $f(n,a) \neq f(m,b)$ :

$$(a1): f(n,a) = n \times (\backslash a^* \backslash + 1) + \#a \neq m \times (\backslash a^* \backslash + 1) + \#a = m \times (\backslash b^* \backslash + 1) + \#b = f(m,b);$$

$$(a2): f(n,a) \bmod a^* = \#a, \text{ whereas } f(m,b) \bmod a^* = f(m,b) \bmod b^* = \#b. \\ \text{and so } f(n,a) \neq f(m,b).^{vi}$$

So in both cases we have:  $\tilde{x} = v_{f(n,a),a^*} \neq v_{f(n,b),a^*} = v_{f(x,b),b^*} = \tilde{y}$ .

ad (b): If  $c \in \mathbf{RT}$  and  $y = v_{m,c}$ , then either (b1)  $m \leq \backslash c \backslash$ , or (b2)  $m > \backslash c \backslash$ . If (b1), then there is exactly one  $a \in [c]$  such that  $m = \#a$  and thus  $f(0,a) = 0 \times (\backslash a^* \backslash + 1) + \#a = m$ , whence  $\widetilde{v_{0,a}} = v_{f(0,a),a^*} = v_{m,c} = y$ . And if (b2), then there is a  $k \leq \backslash c \backslash$  such that  $\backslash c \backslash \bmod m = k$  and thus an  $a \in [c]$  such that  $\#a = k$ ; so if we put  $n := \frac{m-k}{\backslash c \backslash + 1}$ , we have:  $f(n,a) = n \times (\backslash a^* \backslash + 1) + \#a = \frac{m-k}{\backslash c \backslash + 1} \times (\backslash c \backslash + 1) + k = m - k + k = m$ , and thus:  $\widetilde{v_{n,a}} = v_{f(n,a),a^*} = v_{m,c} = y$ .

### A.2.2 Fn. 30

(i): (te)(ee) is logical.

This is because  $id_e$  is ee-invariant, and thus so is  $\{(f, id_e) | f \in D_{te}\} [\in D_{(te)(ee)}]$ , as is readily verified.

(ii):  $D_{(te)(ee)}$  contains an injective function.

In view of the basic set-theoretic fact that, for any  $a, b \in \mathbf{2T}$ ,  $\overline{\overline{D_{ab}}} = \overline{\overline{D_b}}^{\overline{D_a}}$  (where  $\overline{\overline{X}}$  is the cardinality of a set  $X$ ), it must be shown that  $\overline{\overline{D_{te}}} \leq \overline{\overline{D_{ee}}}$ , for any  $n := \overline{\overline{D_e}}$ . Now, if  $n$  is infinite, then  $\overline{\overline{D_{te}}} = n^2 = n < 2^n = n^n = \overline{\overline{D_{ee}}}$ , by basic cardinal arithmetic.<sup>vii</sup> And if  $n$  is finite, 3 cases may be distinguished:

<sup>v</sup>This is one of the standard disjointness assumptions about logical notation that we have been making all along.

<sup>vi</sup>Here  $n \bmod k$  is the remainder of the (Euclidian) division of the dividend  $n$  by the divisor  $k$ .

<sup>vii</sup>See, e.g., Ebbinghaus [2, p. 142, 2.12].

$n = 1$  : Then  $\overline{\overline{D_{te}}} = \overline{\overline{D_e}}^{\overline{\overline{D_t}}} = 1^2 = 1 = 1^1 = \overline{\overline{D_e}}^{\overline{\overline{D_e}}}$ .

$n = 2$  : Then  $\overline{\overline{D_e}} = \overline{\overline{D_t}}$  and thus  $\overline{\overline{D_{te}}} = \overline{\overline{D_{ee}}}$ .

$n > 2$  : Then  $n = m + 2$  (where  $m > 0$ ) and:

$$\overline{\overline{D_{te}}} = \overline{\overline{D_e}}^{\overline{\overline{D_t}}} = n^2 \leq n^m \times n^2 = n^{m+2} = n^n = \overline{\overline{D_{ee}}}.$$

(iii): (te) is not embeddable into (ee).

It suffices to find some  $D_e$  such that  $D_{(te)(ee)}$  does not contain any invariant 1-1 function. So let  $\overline{\overline{D_e}}$  be 2. Then  $\overline{\overline{D_{te}}} = \overline{\overline{D_{ee}}} = 2^2 = 4$  is finite and so any injective  $\mathbf{F} \in D_{(te)(ee)}$  is surjective; in particular,  $id_e = \mathbf{F}(\mathbf{f})$ , for some  $\mathbf{f} \in D_{te}$ . But then, given such  $\mathbf{F}$  and  $\mathbf{f}$  and the permutation  $\pi_e$  that swaps the two members of  $D_e$ , we would also have:

$$\begin{aligned} & \mathbf{F}(\mathbf{f}) \\ = & \pi_{ee}(\mathbf{F}(\mathbf{f})) && \text{invariance of } \mathbf{F}(\mathbf{f}) [= id_e] \\ = & \pi_{(te)(ee)}(\mathbf{F})(\pi_{te}(\mathbf{f})) && \text{definition of } \pi_{(te)(ee)} \\ = & \mathbf{F}(\pi_{te}(\mathbf{f})) && \text{invariance of } \mathbf{F} \end{aligned}$$

Given the injectivity of  $\mathbf{F}$ , it then follows that  $\mathbf{f} = \pi_{te}(\mathbf{f})$ , which cannot be because:

$$\begin{aligned} & \pi_{te}(\mathbf{f})(1) \\ = & \pi_{te}(\mathbf{f})(\pi_t(1)) && \pi_t = id_t \\ = & \pi_e(\mathbf{f}(1)) && \text{definition of } \pi_{(te)(ee)} \\ \neq & \mathbf{f}(1) && \text{definition of } \pi_e \end{aligned}$$

### A.3 Lemmas

#### A.3.1 $\kappa$ -Lemma

Claim:  $\dot{\kappa}_a$  is injective (one-one).

We reason inductively on  $a \in \mathbf{FT}$  and make use of the characterizations of  $\dot{\kappa}_a$  observed in  $\mathbf{k}$ .

If  $\mathbf{f}, \mathbf{f}' \in D_a$  and  $\mathbf{f} \neq \mathbf{f}'$ , where  $a = bc$ , then  $\mathbf{f}(\mathbf{x}_0) \neq \mathbf{f}'(\mathbf{x}_0)$ , for some  $\mathbf{x}_0 \in D_b$ . Hence (!)  $\dot{\kappa}_c(\mathbf{f}(\mathbf{x}_0)) \neq \dot{\kappa}_c(\mathbf{f}'(\mathbf{x}_0))$ , by induction. There are three cases:

- $b = b^*$  :  
then, by (i),  $\dot{\kappa}_c(\mathbf{f}(\mathbf{x}_0)) = \dot{\kappa}_a(\mathbf{f})(\mathbf{x}_0)$  and  $\dot{\kappa}_c(\mathbf{f}'(\mathbf{x}_0)) = \dot{\kappa}_a(\mathbf{f}')(\mathbf{x}_0)$ .  
So, given (!),  $\dot{\kappa}_a(\mathbf{f})(\mathbf{x}_0) \neq \dot{\kappa}_a(\mathbf{f}')(\mathbf{x}_0)$  and thus  $\dot{\kappa}_a(\mathbf{f}) \neq \dot{\kappa}_a(\mathbf{f}')$ .
- $b \neq b^*$  and  $c \in \mathbf{BT}$  :  
then, by (i),  $\dot{\kappa}_c(\mathbf{f}(\mathbf{x}_0)) = \dot{\kappa}_a(\mathbf{f})(\dot{\kappa}_b(\mathbf{x}_0))$  and  $\dot{\kappa}_c(\mathbf{f}'(\mathbf{x}_0)) = \dot{\kappa}_a(\mathbf{f}')(\dot{\kappa}_b(\mathbf{x}_0))$ .  
So, given (!),  $\dot{\kappa}_a(\mathbf{f})(\dot{\kappa}_b(\mathbf{x}_0)) \neq \dot{\kappa}_a(\mathbf{f}')(\dot{\kappa}_b(\mathbf{x}_0))$  and thus  $\dot{\kappa}_a(\mathbf{f}) \neq \dot{\kappa}_a(\mathbf{f}')$ .

- $b \neq b^*$  and  $c \notin \mathbf{BT}$ :  
then, given (!),  $/\dot{\kappa}_c(\mathbf{f}(\mathbf{x}_0))/ \neq / \dot{\kappa}_c(\mathbf{f}'(\mathbf{x}_0))/$ ; but, by (i),  $/\dot{\kappa}_c(\mathbf{f}(\mathbf{x}_0))/ = \dot{\kappa}_a(\mathbf{f})(\dot{\kappa}_b(\mathbf{x}_0))$  and  $/\dot{\kappa}_c(\mathbf{f}'(\mathbf{x}_0))/ = \dot{\kappa}_a(\mathbf{f}')(\dot{\kappa}_b(\mathbf{x}_0))$ . So,  $\dot{\kappa}_a(\mathbf{f})(\dot{\kappa}_b(\mathbf{x}_0)) \neq \dot{\kappa}_a(\mathbf{f}')(\dot{\kappa}_b(\mathbf{x}_0))$  and thus  $\dot{\kappa}_a(\mathbf{f}) \neq \dot{\kappa}_a(\mathbf{f}')$ .

If  $\mathbf{c}, \mathbf{c}' \in D_{se}$ , and  $\mathbf{c} \neq \mathbf{c}'$ , then  $\dot{\kappa}_{se}(\mathbf{c}) = \overleftarrow{\mathbf{c}} \neq \overleftarrow{\mathbf{c}'} = \dot{\kappa}_{se}(\mathbf{c}')$ , as pointed out in fn. 6.

If  $\mathbf{F}, \mathbf{F}' \in D_a$  and  $\mathbf{F} \neq \mathbf{F}'$ , where  $a = s(bc)$  and  $a^* = (a^*(sb^*))$ , then  $\mathbf{F}(\mathbf{w}_0)(\mathbf{x}_0) \neq \mathbf{F}'(\mathbf{w}_0)(\mathbf{x}_0)$ , for some  $\mathbf{w}_0 \in D_s$  and  $\mathbf{x}_0 \in D_b$ ; thus:

$$\begin{aligned} & [\hat{\mathbf{w}} : \mathbf{F}(\mathbf{w})(\mathbf{x}_0)](\mathbf{w}_0) \neq [\hat{\mathbf{w}} : \mathbf{F}'(\mathbf{w})(\mathbf{x}_0)](\mathbf{w}_0), \text{ and so:} \\ & [\hat{\mathbf{w}} : \mathbf{F}(\mathbf{w})(\mathbf{x}_0)] \neq [\hat{\mathbf{w}} : \mathbf{F}'(\mathbf{w})(\mathbf{x}_0)]. \end{aligned}$$

By induction, we thus have:

$$(!) \dot{\kappa}([\hat{\mathbf{w}} : \mathbf{F}(\mathbf{w})(\mathbf{x}_0)]) \neq \dot{\kappa}([\hat{\mathbf{w}} : \mathbf{F}'(\mathbf{w})(\mathbf{x}_0)]).$$

Now, by (i),  $\dot{\kappa}([\hat{\mathbf{w}} : \mathbf{F}(\mathbf{w})(\mathbf{x}_0)]) = \dot{\kappa}(\mathbf{F})(\dot{\kappa}(\mathbf{x}_0))$  and  $\dot{\kappa}([\hat{\mathbf{w}} : \mathbf{F}'(\mathbf{w})(\mathbf{x}_0)]) = \dot{\kappa}(\mathbf{F}')(\dot{\kappa}(\mathbf{x}_0))$ . So, given (!),  $\dot{\kappa}(\mathbf{F})(\dot{\kappa}(\mathbf{x}_0)) \neq \dot{\kappa}(\mathbf{F}')(\dot{\kappa}(\mathbf{x}_0))$  and thus  $\dot{\kappa}(\mathbf{F}) \neq \dot{\kappa}(\mathbf{F}')$ .

If  $\mathbf{C}, \mathbf{C}' \in D_a$  and  $\mathbf{C} \neq \mathbf{C}'$ , where  $a = s(sb)$ , then  $\mathbf{C}(\mathbf{w}_0) \neq \mathbf{C}'(\mathbf{w}_0)$ , for some  $\mathbf{w}_0 \in D_s$ . Hence  $\dot{\kappa}(\mathbf{C}(\mathbf{w}_0)) \neq \dot{\kappa}(\mathbf{C}'(\mathbf{w}_0))$ , by induction, and so:  $\mathbf{c} = \{(\mathbf{w}, (\mathbf{C}(\mathbf{w}))) \mid \mathbf{w} \in D_s\} \neq \{(\mathbf{w}, \dot{\kappa}(\mathbf{C}'(\mathbf{w}))) \mid \mathbf{w} \in D_s\} = \mathbf{c}'$ . But then, by the injectivity of  $\overleftarrow{\cdot}$  (fn. 6):  $\dot{\kappa}(\mathbf{C}) = \overleftarrow{\mathbf{c}} \neq \overleftarrow{\mathbf{c}'} = \dot{\kappa}(\mathbf{C}')$ .

### A.3.2 $\Gamma$ -Lemma

To be shown:  $\dot{\Gamma}_a = "rge(\dot{\kappa}_a)"$ . The proof proceeds by induction on  $a \in \mathbf{FT}$ :<sup>viii</sup> " $\subseteq$ ": Given any  $\mathbf{X} \in D_{a^*}$  such that  $\llbracket \Gamma_a \rrbracket(\mathbf{X}) = 1$  (where  $a \in \mathbf{FT}$ ), we need to find some  $\mathbf{x} \in D_{a^*}$  such that  $\mathbf{X} = \dot{\kappa}_a(\mathbf{x})$ .

– If  $a \in \{t, e, st\}$ ,  $\mathbf{x} = \mathbf{X}$ .

– If  $a = (bc)$ , where  $b = b^*$ , and  $\llbracket \Gamma_a \rrbracket(\mathbf{F}) = 1$ , then for any  $\mathbf{x} \in D_b$ ,  $\llbracket \Gamma_c \rrbracket(\mathbf{F}(\mathbf{x})) = 1$ , and thus by induction,  $\mathbf{F}(\mathbf{x}) = \dot{\kappa}_c(\mathbf{y})$ , for some  $\mathbf{y} = \dot{\kappa}_c^{-1}(\mathbf{F}(\mathbf{x}))$ , given that  $\dot{\kappa}$  is injective. So we put:

$$\mathbf{f} := \{(\mathbf{x}, \dot{\kappa}_c^{-1}(\mathbf{F}(\mathbf{x}))) \mid \mathbf{x} \in D_b\}$$

and observe:

$$\begin{aligned} \dot{\kappa}_a(\mathbf{f}) &= \{(\mathbf{x}, \dot{\kappa}_c(\mathbf{f}(\mathbf{x}))) \mid \mathbf{x} \in D_b\} = \\ & \{(\mathbf{x}, \dot{\kappa}_c(\dot{\kappa}_c^{-1}(\mathbf{F}(\mathbf{x})))) \mid \mathbf{x} \in D_b\} = \{(\mathbf{x}, \mathbf{F}(\mathbf{x})) \mid \mathbf{x} \in D_b\} = \mathbf{F}. \end{aligned}$$

– If  $a = (bc)$ , where  $b \neq b^*$  but  $c \in \mathbf{BT}$ , and  $\llbracket \Gamma_a \rrbracket(\mathbf{F}) = 1$ , then  $\llbracket \Gamma_c \rrbracket(\mathbf{F}(\mathbf{X})) = 1$ , for any  $\mathbf{X}$  such that  $\llbracket \Gamma_b \rrbracket(\mathbf{X}) = 1$ ; hence by induction,  $\mathbf{F}(\mathbf{X}) \in rge(\dot{\kappa}_c)$  whenever  $\mathbf{X} \in rge(\dot{\kappa}_b)$ . We now put:

$$\mathbf{f} := \{(\dot{\kappa}_b^{-1}(\mathbf{X}), \dot{\kappa}_c^{-1}(\mathbf{F}(\mathbf{X}))) \mid \mathbf{X} \in rge(\dot{\kappa}_b)\}.$$

Clearly,  $\mathbf{f} \in D_a$ : any  $\mathbf{x} \in D_b$  is of the form  $\dot{\kappa}_b^{-1}(\mathbf{X})$ , where  $\mathbf{X} = \dot{\kappa}_b(\mathbf{x}) \in$

<sup>viii</sup>Though the presentation separates the directions, the induction is understood as treating them simultaneously, so that at each step the inductive hypothesis has access to the full biconditional.

$rge(\dot{\kappa}_b)$ ; and  ${}^{-1}_c(\mathbf{F}(\mathbf{X})) \in D_c$ . It remains to be shown that  $\dot{\kappa}_a(\mathbf{f}) = \mathbf{F}$ , i.e., that for any  $\mathbf{X} \in D_{b^*}$ :

$$(!) \quad \dot{\kappa}_a(\mathbf{f})(\mathbf{X}) = \mathbf{F}(\mathbf{X}).$$

Two cases must be distinguished, according to whether  $\mathbf{X} \in rge(\dot{\kappa}_b)$ . If so,  $\mathbf{X} = \dot{\kappa}_b(\mathbf{x})$ , for some  $\mathbf{x} \in D_b$ , and we may conclude:

$$\begin{aligned} {}_a(\mathbf{f})(\mathbf{X}) &= {}_a(\mathbf{f})(\dot{\kappa}_b(\mathbf{x})) = {}_c(\mathbf{f}(\mathbf{x})) = {}_c(\mathbf{f}(\dot{\kappa}_b^{-1}(\mathbf{X}))) = \\ &\dot{\kappa}_c(\dot{\kappa}_c^{-1}(\mathbf{F}(\mathbf{X}))) = \mathbf{F}(\mathbf{X}), \end{aligned}$$

thus establishing the first case of (!). But if  $\mathbf{X} \notin rge(\dot{\kappa}_b)$ , then by induction,  $\llbracket \Gamma_b \rrbracket(\mathbf{X}) = 0$ , and so,<sup>ix</sup> by the definitions of  $\Gamma_a$  and  $\dot{\kappa}_a$ :  $\mathbf{F}(\mathbf{X}) = \mathbf{0}_{c^*} = \dot{\kappa}_a(\mathbf{f})(\mathbf{X})$ .

– If  $a = (bc)$ ,  $b \neq b^*$ ,  $c \notin \mathbf{BT}$ , and  $\llbracket \Gamma_a \rrbracket(\mathbf{F}) = 1$ , then by induction, for any  $\mathbf{X} \in rge(\dot{\kappa}_b)$ , there is a  $\mathbf{Y} \in rge(\dot{\kappa}_c)$  such that  $\mathbf{F}(\mathbf{X}) = \mathbf{Y}/$ . We now put:

$$\mathbf{f} := \{(\dot{\kappa}_b^{-1}(\mathbf{X}), \mathbf{y}) \mid \mathbf{F}(\mathbf{X})(\dot{\kappa}_c(\mathbf{y})) = 1\}.$$

Clearly,  $\mathbf{f} \in D_a$ : any  $\mathbf{x} \in D_b$  is of the form  $\dot{\kappa}_b^{-1}(\mathbf{X})$ , where  $\mathbf{X} = \dot{\kappa}_b(\mathbf{x}) \in rge(\dot{\kappa}_b)$ ; and given that  $\mathbf{F} \in D_a$ ,  $\mathbf{y} \in D_c$  if  $\mathbf{F}(\mathbf{X})(\dot{\kappa}_c(\mathbf{y})) = 1$ .<sup>x</sup> Hence, for any  $\mathbf{X} \in rge(\dot{\kappa}_b)$  and  $\mathbf{Y} \in D_{c^*}$ :

$$\begin{aligned} &\dot{\kappa}_a(\mathbf{f})(\mathbf{X})(\mathbf{Y}) = 1 \\ \text{iff} \quad &\mathbf{Y} = \dot{\kappa}_c(\mathbf{f}(\dot{\kappa}_b^{-1}(\mathbf{X}))) && \text{by the definition of } \dot{\kappa}_c \\ \text{iff} \quad &\dot{\kappa}_c^{-1}(\mathbf{Y}) = \mathbf{f}(\dot{\kappa}_b^{-1}(\mathbf{X})) && \text{by the injectivity of } \dot{\kappa}_c \\ \text{iff} \quad &\mathbf{F}(\mathbf{X})(\dot{\kappa}_c(\dot{\kappa}_c^{-1}(\mathbf{Y}))) = 1 && \text{by the definition of } \mathbf{f} \\ \text{iff} \quad &\mathbf{F}(\mathbf{X})(\mathbf{Y}) = 1 && \text{by the injectivity of } \dot{\kappa}_c \end{aligned}$$

On the other hand, if  $\mathbf{F} \notin rge(\dot{\kappa}_b)$ , then  $\llbracket \Gamma_b \rrbracket(\mathbf{X}) = 0$ , by induction. Hence, as in the case before, the definitions of  $\Gamma_a$  and  $\dot{\kappa}_a$  guarantee that  $\mathbf{F}(\mathbf{X}) = \mathbf{0}_{c^*} = \dot{\kappa}_a(\mathbf{f})(\mathbf{X})$ .

– If  $a = (\text{se})$  and  $\llbracket \Gamma_a \rrbracket(\mathbf{R}) = 1$ , then for any  $\mathbf{w} \in D_s$  there is a unique  $\mathbf{u} \in D_e$  such that  $\mathbf{R}(\mathbf{u})(\mathbf{w}) = 1$ . So we put:

$$\mathbf{c} := \{(\mathbf{w}, \mathbf{u}) \in D_s \times D_e \mid \mathbf{R}(\mathbf{u})(\mathbf{w}) = 1\} \in D_{\text{se}},$$

and we have, for any  $\mathbf{u}_0 \in D_e$  and  $\mathbf{w}_0 \in D_s$ :

$$\begin{aligned} &{}_a(\mathbf{c})(\mathbf{u}_0)(\mathbf{w}_0) = 1 \\ \text{iff} \quad &\overleftarrow{\mathbf{c}}(\mathbf{u}_0)(\mathbf{w}_0) = 1 && \text{by the definition of } \dot{\kappa}_a \\ \text{iff} \quad &\mathbf{c}(\mathbf{w}_0) = \mathbf{u}_0 && \text{cf. fn. 6} \\ \text{iff} \quad &\mathbf{R}(\mathbf{u}_0)(\mathbf{w}_0) = 1. && \text{by the definition of } \mathbf{c} \end{aligned}$$

Hence  $\mathbf{R} = \dot{\kappa}_a(\mathbf{c})$ , as required.

– If  $a = (\text{s}(bc))$ , and  $\llbracket \Gamma_a \rrbracket(\mathbf{H}) = 1$ , then  $\llbracket \Gamma_{sc} \rrbracket(\mathbf{H}(\mathbf{X})) = 1$ , for any  $\mathbf{X}$  such that  $\llbracket \Gamma_b \rrbracket(\mathbf{X}) = 1$ . Hence by induction,  $\mathbf{H}(\mathbf{X}) \in rge(\dot{\kappa}_{sc})$  whenever  $\mathbf{X} \in rge(\dot{\kappa}_b)$ . We now put:

$$\mathbf{F} := [\hat{\mathbf{w}} : \{(\dot{\kappa}_b^{-1}(\mathbf{X}), \dot{\kappa}_{sc}^{-1}(\mathbf{H}(\mathbf{X}))(\mathbf{w})) \mid \mathbf{X} \in rge(\dot{\kappa}_b)\}] \quad [\in D_a]$$

and it remains to be shown that:

$$(!!) \quad \dot{\kappa}_a(\mathbf{F})(\mathbf{X}) = \mathbf{H}(\mathbf{X}),$$

for any  $\mathbf{X} \in D_{b^*}$ . Now, two cases need to be distinguished:

<sup>ix</sup>Here the architectural remarks in fn. viii come to bear.

<sup>x</sup>The argument capitalizes on the simplification addressed in fn. 5.

Case 1: for some  $x \in D_b$ ,  $X = \dot{\kappa}_b(x)$ ; we then observe:

$$\begin{aligned}
& \dot{\kappa}_a(\mathbf{F})(X) \\
= & {}_a(\mathbf{F})(\dot{\kappa}_b(x)) & X = \dot{\kappa}_b(x) \\
= & {}_{sc}([\hat{w} : \mathbf{F}(w)(x)]) & \text{by the definition of } \dot{\kappa}_a \\
= & \dot{\kappa}_{sc}([\hat{w} : \dot{\kappa}_{sc}^{-1}(\mathbf{H}(X))(w)]) & \mathbf{F}(w)(x) = {}_{sc}^{-1}(\mathbf{H}(X))(w) \\
= & \dot{\kappa}_{sc}(\dot{\kappa}_{sc}^{-1}(\mathbf{H}(X))) & \text{by } \eta\text{-conversion}^{\text{xi}} \\
= & \mathbf{H}(X) & \dot{\kappa}_{sc} \circ \dot{\kappa}_{sc}^{-1} = \text{id}_{sc}
\end{aligned}$$

Case 2:  $X \notin \text{rge}(\dot{\kappa}_b)$ ,  $\dot{\kappa}_a(\mathbf{F})(X) = \mathbf{0}_{(sc)^*}$ , by the definition of  $\dot{\kappa}_a$ , and  $\mathbf{H}(X) = \mathbf{0}_{(sc)^*}$ , by the definition of  $\Gamma_a$  and the inductive hypothesis.

– If  $a = (s(sb))$ , where  $b \in \mathbf{FT}$ , and  $\llbracket \Gamma_a \rrbracket(\mathbf{R}) = 1$ , then by induction, for any  $w \in D_s$ , there is a unique  $X_w \in D_{(sb)^*}$  such that  $\mathbf{R}(X_w)(w) = 1$ ; moreover,  $X_w \in \text{rge}(\dot{\kappa}_{sb})$ . We now put:

$$C := [\hat{w}_1 : [\hat{w}_2 : \dot{\kappa}_{sb}^{-1}(X_{w_1})(w_2)]]$$

Clearly,  $C \in D_a$ , and we need to show that,  $\dot{\kappa}_a(C) = \mathbf{R}$ . And in fact, for any  $Y \in D_{(sb)^*}$  and  $w_0 \in D_s$  we have:

$$\begin{aligned}
& \dot{\kappa}_a(C)(Y)(w_0) = 1 \\
\text{iff } & [\hat{w} : \dot{\kappa}_{sb}(C(w))](Y)(w_0) = 1 & \text{by the definition of } \dot{\kappa}_a \\
\text{iff } & [\hat{w} : \dot{\kappa}_{sb}(C(w))](w_0) = Y & \text{cf. fn 6} \\
\text{iff } & \dot{\kappa}_{sb}(C(w_0)) = Y & \text{by } \beta\text{-conversion}^{\text{xii}} \\
\text{iff } & \dot{\kappa}_{sb}([\hat{w}_1 : [\hat{w}_2 : \dot{\kappa}_{(sb)^*}^{-1}(X_{w_1})(w_2)]](w_0)) = Y & \text{by the definition of } C \\
\text{iff } & \dot{\kappa}_{sb}([\hat{w}_2 : \dot{\kappa}_{sb}^{-1}(X_{w_0})(w_2)]) = Y & \text{by } \beta\text{-conversion} \\
\text{iff } & \dot{\kappa}_{sb}(\dot{\kappa}_{sb}^{-1}(X_{w_0})) = Y & \text{by } \eta\text{-conversion} \\
\text{iff } & X_{w_0} = Y & \dot{\kappa}_{sb} \circ \dot{\kappa}_{sb}^{-1} = \text{id}_b \\
\text{iff } & \mathbf{R}(Y)(w_0) = 1. & \text{by the definition of } X_{w_0}
\end{aligned}$$

“ $\supseteq$ ”: We need to show that, for any  $x \in D_a$ ,  $\llbracket \Gamma_a \rrbracket(\dot{\kappa}_a(x)) = 1$ . As in the forward direction, the base cases  $a \in \{t, e, st\}$  are trivial.

– If  $a = (bc)$ , where  $b = b^*$ , and  $f \in D_a$ , we must show that  $\llbracket \Gamma_a \rrbracket(\dot{\kappa}_a(f)) = 1$ , which, by the definition of  $\Gamma_a$ , means that for any  $x \in D_b$ ,  $\llbracket \Gamma_c \rrbracket(\dot{\kappa}_a(f)(x)) = 1$ . By induction, this in turn means that  ${}_a(f)(x) \in \text{rge}(\dot{\kappa}_c)$ . But then by the definition of  $\dot{\kappa}_a : \dot{\kappa}_a(f)(x) = \dot{\kappa}_c(f(x))$ , and so  $\dot{\kappa}_a(f)(x) \in \text{rge}(\dot{\kappa}_c)$ , indeed.

– If  $a = (bc)$ , where  $b \neq b^*$  but  $c \in \mathbf{BT}$ , and  $f \in D_a$ , it must be shown that, for any  $X \in D_{b^*}$ : (1)  $\llbracket \Gamma_c \rrbracket(\dot{\kappa}_a(f)(X)) = 1$  if  $\llbracket \Gamma_b \rrbracket(X) = 1$ ; and that (2)  $\dot{\kappa}_a(f)(X) = \mathbf{0}_{c^*}$  if  $\llbracket \Gamma_b \rrbracket(X) = 0$ . By induction, (1) means that  $\dot{\kappa}_a(f)(X) \in \text{rge}(\dot{\kappa}_c)$  if  $X = \dot{\kappa}_b(x)$ , for some  $x \in D_b$ ; but then, by the definition of  $\dot{\kappa}_a : \dot{\kappa}_a(f(\dot{\kappa}_b(x))) = \dot{\kappa}_c(f(X)) \in \text{rge}(\dot{\kappa}_c)$ . And (2) holds because by induction,  $\llbracket \Gamma_b \rrbracket(X) = 0$  implies:  $X \notin \text{rge}(\dot{\kappa}_b)$ , and thus by the definition of  $\dot{\kappa}_a : \dot{\kappa}_a(f)(X) = \mathbf{0}_{c^*}$ , indeed.

<sup>xi</sup>The functions  $f$  denoted by our meta-linguistic abstraction operator “ $\hat{w}$ ” are extensional and thus satisfy the equation:  $[\hat{w} : f(w)] = f$ , aka  $\eta$ -conversion.

<sup>xii</sup>This is the logical principle, according to which  $[\hat{w} : \dots w \dots](w_0) = \dots w_0 \dots$ , provided that  $w_0$  does not get bound in the substitution process.

– If  $a = (bc)$ ,  $b \neq b^*$ ,  $c \notin \mathbf{BT}$ , and  $\mathbf{f} \in D_a$ , it must be shown that (1) for any  $\mathbf{X} \in rge(\dot{\kappa}_b) : \dot{\kappa}_a(\mathbf{f})(\mathbf{X}) = / \mathbf{Y} /$ , for some  $\mathbf{Y} \in rge(\dot{\kappa}_c)$ ; and that (2)  $\dot{\kappa}_a(\mathbf{f})(\mathbf{X}) = \mathbf{0}_{c^*t}$  whenever  $\mathbf{X} \in D_{b^*} \setminus rge(\dot{\kappa}_b)$ . (2) follows as in the previous case. In case (1),  $\mathbf{X} = \dot{\kappa}_b(\mathbf{X})$  again, whence,  $\dot{\kappa}_a$ , we have, following the definition of  $\dot{\kappa}_a : \dot{\kappa}_a(\mathbf{f})(\mathbf{X}) = / \dot{\kappa}_c(\mathbf{f}(\mathbf{x})) /$ , which is of the required form.

– If  $a = (se)$  and  $\mathbf{c} \in D_a$ , we must show that, for any  $\mathbf{w} \in D_w$ , there is exactly one  $\mathbf{u}_w \in D_e$  such that:  $\dot{\kappa}_a(\mathbf{c})(\mathbf{u}_w)(\mathbf{w}) = 1$ . If we put:  $\mathbf{u}_w = \mathbf{c}(\mathbf{w})$ , the definition of  $\dot{\kappa}_a$  yields.  $\dot{\kappa}_a(\mathbf{c})(\mathbf{u}_w)(\mathbf{w}) = 1$  iff  $\mathbf{c}(\mathbf{w}) = \mathbf{u}_w$ , which is how we just defined  $\mathbf{u}_w$ ; and its uniqueness follows because, for any  $\mathbf{v} \neq \mathbf{u}_w$ , we get:  $\dot{\kappa}_a(\mathbf{c})(\mathbf{v})(\mathbf{w}) = 1$  iff  $\mathbf{c}(\mathbf{w}) = \mathbf{v}$ , which is not the case, by assumption. – If  $a = (s(bc))$  and  $\mathbf{F} \in D_a$ , it must be shown that (1) for any  $\mathbf{X} \in rge(\dot{\kappa}_b) : \dot{\kappa}_a(\mathbf{F})(\mathbf{X}) \in rge(\dot{\kappa}_{sb})$ ; and that (2)  $\dot{\kappa}_a(\mathbf{F})(\mathbf{X}) = \mathbf{0}_{(sc)^*}$  whenever  $\mathbf{X} \in D_{b^*} \setminus rge(\dot{\kappa}_b)$ . Again, (1) holds because  $\mathbf{X} = \dot{\kappa}_b(\mathbf{x})$ , where  $\mathbf{x} \in D_b$ , and thus  $\dot{\kappa}_a(\mathbf{F})(\mathbf{X}) = \dot{\kappa}_{sc}([\hat{\mathbf{w}} : \mathbf{F}(\mathbf{w})(\mathbf{x})])$ . And (2) follows as in previous cases.

– If  $a = (s(sb))$  and  $\mathbf{C} \in D_a$ , we need to show that, for any  $\mathbf{w} \in D_s$ , there is a unique  $\mathbf{X}_w \in D_{(sb)^*}$  such that  $\dot{\kappa}_a(\mathbf{C})(\mathbf{X}_w)(\mathbf{w}) = 1$  where  $\mathbf{X}_w \in rge(\dot{\kappa}_{sb})$ . We put:  $\mathbf{X}_w := \dot{\kappa}_{sb}(\mathbf{C}(\mathbf{w}))$  and observe that  $\mathbf{X}_w \in rge(\dot{\kappa}_{sb})$  and that:

$$\begin{aligned} & \dot{\kappa}_a(\mathbf{C})(\mathbf{X}_w)(\mathbf{w}) = 1 \\ \text{iff } & [\hat{\mathbf{w}}' : \dot{\kappa}_{sb}(\mathbf{C}(\mathbf{w}'))](\mathbf{X}_w)(\mathbf{w}) = 1 \quad \text{by the definition of } \dot{\kappa}_a \text{ (and } \alpha\text{-conversion)}^{\text{xiii}} \\ \text{iff } & [\hat{\mathbf{w}}' : \dot{\kappa}_{sb}(\mathbf{C}(\mathbf{w}'))](\mathbf{w}) = \mathbf{X}_w \quad \text{cf. fn. 6} \\ \text{iff } & \dot{\kappa}_{sb}(\mathbf{C}(\mathbf{w})) = \mathbf{X}_w \quad \text{by } \beta\text{-conversion,} \end{aligned}$$

which is how  $\mathbf{X}_w$  was defined. And again, its uniqueness follows as a corollary of the above calculation, trading  $\mathbf{X}_w$  for arbitrary  $\mathbf{Y} \in D_{(sb)^*}$ .

### A.3.3 $\Phi$ -Lemma

To be shown:

$$\begin{aligned} & \dot{\Phi}_{b,c}(\mathbf{F})(\mathbf{X})(\mathbf{Y}) = 1 \text{ iff} \\ & \text{there are } \mathbf{f} \in D_{bc}, \mathbf{x} \in D_b, \text{ and } \mathbf{y} \in D_c \text{ such that:} \\ & \quad \mathbf{F} = \dot{\kappa}_{ab}(\mathbf{f}), \mathbf{X} = \dot{\kappa}_b(\mathbf{x}), \mathbf{Y} = \dot{\kappa}_c(\mathbf{y}), \text{ and } \mathbf{f}(\mathbf{x}) = \mathbf{y}. \end{aligned}$$

If  $\dot{\Phi}_{b,c}(\mathbf{F})(\mathbf{X})(\mathbf{Y}) = 1$ , two cases are to be distinguished. First, if  $(bc)^* = b^*c^*$ , then  $\dot{\Gamma}_{bc}(\mathbf{F}) = \dot{\Gamma}_b(\mathbf{X}) = 1$  and  $\mathbf{F}(\mathbf{X}) = \mathbf{Y}$ , by the interpretation of  $IL^*$ . Hence  $\mathbf{F} \in rge(\dot{\kappa}_{bc})$  and  $\mathbf{X} \in rge(\dot{\kappa}_b)$ , by the  $\Gamma$ -Lemma. So there are  $\mathbf{f} \in D_{bc}$  and  $\mathbf{x} \in D_b$  such that  $\mathbf{F} = \dot{\kappa}_{bc}(\mathbf{f})$  and  $\mathbf{X} = \dot{\kappa}_b(\mathbf{x})$ , as required. Moreover,  $\mathbf{Y} = \mathbf{F}(\mathbf{X}) = \dot{\kappa}_{bc}(\mathbf{f})(\dot{\kappa}_b(\mathbf{x})) = \dot{\kappa}_c(\mathbf{f}(\mathbf{x}))$ , by the definition of  $\dot{\kappa}_{bc}$ , and thus the lemma holds for  $\mathbf{y} := \mathbf{f}(\mathbf{x})$ . The second case is similar:  $\mathbf{f}$  and  $\mathbf{x}$  exist by the  $\Gamma$ -Lemma, and if we put  $\mathbf{y} := \mathbf{f}(\mathbf{x})$  again, the definition of  $\dot{\kappa}_{bc}$  gives us:  $\mathbf{F}(\mathbf{X}) = \dot{\kappa}_{bc}(\mathbf{f})(\dot{\kappa}_b(\mathbf{x})) = / \dot{\kappa}_c(\mathbf{f}(\mathbf{x})) / = / \dot{\kappa}_c(\mathbf{y}) / = / \mathbf{Y} /$ , by the third conjunct of  $\dot{\Phi}_{bc}$  – and thus:  $\dot{\kappa}_c(\mathbf{y}) = \mathbf{Y}$ , since  $/ \cdot /$  is (obviously) injective.

<sup>xiii</sup>In other words, the variable bound by the functional abstraction “ $\hat{\mathbf{w}}$ ” may be renamed, if this causes no confusion.

### A.3.4 $\Psi$ -Lemma

To be shown:  $\dot{\Psi}_a(\mathbf{C})(\mathbf{X})(\mathbf{w}) = 1$  iff

there are  $\mathbf{c} \in D_{sa}$  and  $\mathbf{x} \in D_a$  such that:

(A)  $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$ , (B)  $\mathbf{X} = \dot{\kappa}_a(\mathbf{x})$ , and (C)  $\mathbf{c}(\mathbf{w}) = \mathbf{x}$ .

We proceed by induction on  $|b|$ .

– If  $a = t$ ,  $\mathbf{p} \in D_{st}$ ,  $\mathbf{v} \in D_t$  and  $\mathbf{w}_0 \in D_s$ , then, by the interpretation of  $IL^*$ :  $\dot{\Psi}_b(\mathbf{p})(\mathbf{v})(\mathbf{w}_0) = 1$  just in case (C)  $\mathbf{p}(\mathbf{w}_0) = \mathbf{v}$ , where (A)  $\dot{\kappa}_{st}(\mathbf{p}) = \mathbf{p}$  and (B)  $\dot{\kappa}_t(\mathbf{v}) = \mathbf{v}$ , due to the characterization **k**) of  $\dot{\kappa}_t$  and  $\dot{\kappa}_{st}$ .

– If  $a = e$ ,  $\mathbf{R} \in D_{e(st)}$ ,  $\mathbf{u} \in D_e$  and  $\mathbf{w}_0 \in D_s$ , then, by the interpretation of  $IL^*$ :  $\dot{\Psi}_b(\mathbf{R})(\mathbf{u})(\mathbf{w}_0) = 1$  just in case  $\dot{\Gamma}(\mathbf{R})(\mathbf{w}_0) = \mathbf{R}(\mathbf{u})(\mathbf{w}_0) = 1$ , which by the  $\Gamma$ -Lemma, means that, for some  $\mathbf{c} \in D_{se}$ :  $\mathbf{R} = \dot{\kappa}_{se}(\mathbf{c})$  and  $\mathbf{R}(\mathbf{u})(\mathbf{w}_0) = 1$ . In other words:  $\dot{\Psi}_b(\mathbf{R})(\mathbf{u})(\mathbf{w}_0) = 1$  iff there is a  $\mathbf{c} \in D_{se}$  such that: (A)  $\mathbf{R} = \dot{\kappa}_{se}(\mathbf{c})$ , (B)  $\mathbf{u} = \dot{\kappa}_e(\mathbf{u})$ , and  $\dot{\kappa}_{(se)}(\mathbf{c})(\dot{\kappa}_e(\mathbf{u}))(\mathbf{w}_0) = 1$ , i.e.: (C)  $\mathbf{c}(\mathbf{w}_0) = \mathbf{u}$ , by the characterization **l**) of  $\dot{\kappa}_{se}$ .

– If  $a = bc$ ,  $\mathbf{G} \in D_{(s(bc))^*}$ ,  $\mathbf{F} \in D_{(bc)^*} = D_{b^*c^*}$  and  $\mathbf{w}_0 \in D_s$ , then  $(s(bc))^* = b^*(sc)^*$ , and:

$\dot{\Psi}_{bc}(\mathbf{G})(\mathbf{F})(\mathbf{w}_0) = 1$   
iff  $\dot{\Gamma}_{s(bc)}(\mathbf{G}) = \dot{\Gamma}_{bc}(\mathbf{F}) = \dot{\Psi}_c(\mathbf{G}(\mathbf{X}))(F(\mathbf{X}))(\mathbf{w}_0) = 1$   
whenever  $\dot{\Gamma}_b(\mathbf{X}) = 1$  interpretation of  $IL^*$   
iff  $\mathbf{G} \in rge(\dot{\kappa}_{s(bc)}), \mathbf{F} \in rge(\dot{\kappa}_{bc})$ , and  $\dot{\Psi}_c(\mathbf{G}(\mathbf{X}))(F(\mathbf{X}))(\mathbf{w}_0) = 1$   
whenever  $\mathbf{X} \in rge(\dot{\kappa}_b)$   $\Gamma$ -Lemma  
iff there are  $\mathbf{g} \in D_{s(bc)}$  and  $\mathbf{f} \in D_{bc}$  such that:  $\dot{\kappa}_{s(bc)}(\mathbf{g}) = \mathbf{G}$ ,  $\dot{\kappa}_{bc}(\mathbf{f}) = \mathbf{F}$ ,  
and for any  $\mathbf{x} \in D_b$  there are  $\mathbf{z} \in D_{sc}$  and  $\mathbf{y} \in D_c$  such that:  
 $\dot{\kappa}_{sc}(\mathbf{z}) = \mathbf{G}(\dot{\kappa}_b(\mathbf{x}))$ ,  $\dot{\kappa}_c(\mathbf{y}) = \mathbf{F}(\dot{\kappa}_b(\mathbf{x}))$ , and  $\mathbf{y} = \mathbf{z}(\mathbf{w}_0)$  induction  
iff there are  $\mathbf{g} \in D_{s(bc)}$  and  $\mathbf{f} \in D_{bc}$  such that:  $\dot{\kappa}_{s(bc)}(\mathbf{g}) = \mathbf{G}$ ,  $\dot{\kappa}_{bc}(\mathbf{f}) = \mathbf{F}$ ,  
and for any  $\mathbf{x} \in D_b$  there is a  $\mathbf{z} \in D_{sc}$  such that:  
 $\dot{\kappa}_{sc}(\mathbf{z}) = \dot{\kappa}_{s(bc)}(\mathbf{g})(\dot{\kappa}_b(\mathbf{x}))$  and  $\dot{\kappa}_c(\mathbf{z}(\mathbf{w}_0)) = \dot{\kappa}_{bc}(\mathbf{f})(\dot{\kappa}_b(\mathbf{x}))$  identity  
iff there are  $\mathbf{g} \in D_{s(bc)}$  and  $\mathbf{f} \in D_{bc}$  such that:  $\dot{\kappa}_{s(bc)}(\mathbf{g}) = \mathbf{G}$ ,  $\dot{\kappa}_{bc}(\mathbf{f}) = \mathbf{F}$ ,  
and for any  $\mathbf{x} \in D_b$  there is a  $\mathbf{z} \in D_{sc}$  such that:  
 $\dot{\kappa}_{sc}(\mathbf{z}) = \dot{\kappa}_{sc}([\hat{\mathbf{w}} : \mathbf{g}(\mathbf{w})(\mathbf{x})])$  and  $\mathbf{f}(\mathbf{x}) = \mathbf{z}(\mathbf{w}_0)$  iff def.  $\dot{\kappa}_{s(bc)}, \dot{\kappa}_{bc}$   
iff there are  $\mathbf{g} \in D_{s(bc)}$  and  $\mathbf{f} \in D_{bc}$  such that:  $\dot{\kappa}_{s(bc)}(\mathbf{g}) = \mathbf{G}$ ,  $\dot{\kappa}_{bc}(\mathbf{f}) = \mathbf{F}$ ,  
and for any  $\mathbf{x} \in D_b$  there is a  $\mathbf{z} \in D_{sc}$  such that:  
 $\mathbf{z} = [\hat{\mathbf{w}} : \mathbf{g}(\mathbf{w})(\mathbf{x})]$  and  $\mathbf{f}(\mathbf{x}) = \mathbf{z}(\mathbf{w}_0)$  injectivity of  $\dot{\kappa}_{sc}$   
iff there are  $\mathbf{g} \in D_{s(bc)}$  and  $\mathbf{f} \in D_{bc}$  such that:  $\dot{\kappa}_{s(bc)}(\mathbf{g}) = \mathbf{G}$ ,  $\dot{\kappa}_{bc}(\mathbf{f}) = \mathbf{F}$ ,  
and for any  $\mathbf{x} \in D_b$ :  $\mathbf{f}(\mathbf{x}) = \mathbf{g}(\mathbf{w}_0)(\mathbf{x})$  identity +  $\beta$ -conversion  
iff there are  $\mathbf{g} \in D_{s(bc)}$  and  $\mathbf{f} \in D_{bc}$  such that:  
(A)  $\dot{\kappa}_{s(bc)}(\mathbf{g}) = \mathbf{G}$ , (B)  $\dot{\kappa}_{bc}(\mathbf{f}) = \mathbf{F}$ , and (C)  $\mathbf{f} = \mathbf{g}(\mathbf{w}_0)$  extensionality

– If  $a = sb$ ,  $G \in D_{(sb)^*(st)}$ ,  $C \in D_{(sb)^*}$  and  $w_0 \in D_s$ , then:

$\dot{\Psi}_{sb}(G)(C)(w_0) = 1$

iff  $\dot{\Gamma}_{s(sb)}(G) = \dot{\Gamma}_{sb}(C) = G(C)(w_0) = 1$  interpretation of  $IL^*$

iff there are  $g \in D_{s(sb)}$  and  $c \in D_{sb}$  such that  $G = \dot{\kappa}_{s(sb)}(g)$ ,  $C = \dot{\kappa}_{sb}(c)$  Γ-Lemma  
and  $G(C)(w_0) = 1$

iff there are  $g \in D_{s(sb)}$  and  $c \in D_{sb}$  such that  $G = \dot{\kappa}_{s(sb)}(g)$ ,  $C = \dot{\kappa}_{sb}(c)$  identity  
and  $\dot{\kappa}_{s(sb)}(g)(\dot{\kappa}_{sb}(c))(w_0) = 1$

iff there are  $g \in D_{s(sb)}$  and  $c \in D_{sb}$  such that  $G = \dot{\kappa}_{s(sb)}(g)$ ,  $C = \dot{\kappa}_{sb}(c)$  def.  $\dot{\kappa}_{s(sb)}$   
and  $[\hat{w} : \dot{\kappa}_{sb}(g(w))](\dot{\kappa}_{sb}(c))(w_0) = 1$

iff there are  $g \in D_{s(sb)}$  and  $c \in D_{sb}$  such that  $G = \dot{\kappa}_{s(sb)}(g)$ ,  $C = \dot{\kappa}_{sb}(c)$  e)  
and  $[\hat{w} : \dot{\kappa}_{sb}(g(w))](w_0) = \dot{\kappa}_{sb}(c)$

iff there are  $g \in D_{s(sb)}$  and  $c \in D_{sb}$  such that  $G = \dot{\kappa}_{s(sb)}(g)$ ,  $C = \dot{\kappa}_{sb}(c)$  β-conversion  
and  $\dot{\kappa}_{sb}(g(w_0)) = \dot{\kappa}_{sb}(c)$

iff there are  $g \in D_{s(sb)}$  and  $c \in D_{sb}$  such that:

(A)  $G = \dot{\kappa}_{s(sb)}(g)$ , (B)  $C = \dot{\kappa}_{sb}(c)$  and (C)  $g(w_0) = c$  injectivity of  $\dot{\kappa}_{sb}$

#### A.4 ω-Lemma

To be shown:

$$(\Omega) \quad \omega(a^*) \leq \omega(a^+).$$

If  $a = t$  or  $a = e$ ,  $a^* = a = a^+$  and thus  $(\Omega)$  holds trivially. – If  $a = (bc)$  (where  $b \neq s$ ), then either (A)  $a^+ = (b^+c^+)$  or else: (B)  $a^+ = (b^+(c^+t))$ . In case (A),  $b = b^+$  or  $c^+ \in \mathbf{BT}$ , and thus:  $b = b^*$  or  $c^* \in \mathbf{BT}$ , by **z**). So  $a^* = (b^*c^*)$  and we have:  $\omega(a^+) = \max(\omega(b^+) + 1, \omega(c^+))$  and  $\omega(a^*) = \max(\omega(b^*) + 1, \omega(c^*))$ . But then  $\omega(a^*) \leq \omega(a^+)$  because, by induction:  $\omega(c^*) \leq \omega(c^+)$  and  $\omega(b^*) \leq \omega(b^+)$  and so  $\omega(b^*) + 1 \leq \omega(b^+) + 1$ . In Case (B),  $a^* = (b^*(c^*t))$ , again by **z**), and so we have:

$$\begin{aligned} & \omega(a^+) \\ = & \max(\omega(b^+) + 1, \max(\omega(c^+) + 1, \omega(t))) \\ = & \max(\omega(b^+) + 1, \max(\omega(c^+) + 1, 0)) \\ = & \max(\omega(b^+) + 1, \omega(c^+) + 1); \end{aligned}$$

and similarly:

$$\omega(a^*) = \max(\omega(b^*) + 1, \omega(c^*) + 1).$$

But then, again,  $\omega(a^*) \leq \omega(a^+)$  because, by induction:  $\omega(b^*) \leq \omega(b^+)$  and  $\omega(c^*) \leq \omega(c^+)$  and so  $\omega(b^*) + 1 \leq \omega(b^+) + 1$  and  $\omega(c^*) + 1 \leq \omega(c^+) + 1$ . – If  $a = (st)$  or  $a = (se)$ ,  $a^* = a = a^+$  and thus  $(\Omega)$  holds trivially. – If  $a = s(bc)$  (where  $b \neq s$ ), then:

$$\begin{aligned} & \omega(a^*) \\ = & \omega(b^*(sc)^*) \\ = & \max(\omega(b^*) + 1, \omega((sc)^*)) \\ \leq & \max(\omega(b^*) + 1, \omega((sc)^+)) & \text{by induction} \\ = & \max(\omega(b^*) + 1, \max(\omega((c)^+) + 1, \omega(st))) & \text{def. } \omega((sc)^+) \\ = & \max(\omega(b^*) + 1, \omega((c)^+) + 1) \end{aligned}$$

Moreover:

$$\begin{aligned}
& \omega(a^+) \\
= & \omega((bc)^+(st)) \\
= & \max(\omega((bc)^+) + 1, \max(\omega(s) + 1, \omega(t))) \\
= & \max(\omega((bc)^+) + 1, 1) \\
= & \omega((bc)^+) + 1
\end{aligned}$$

It thus remains to be shown that  $(+) \max(\omega(b^+) + 1, \omega(c^+) + 1) \leq \omega((bc)^+) + 1$ , i.e., that  $\max(\omega(b^+), \omega(c^+)) \leq \omega((bc)^+)$ . There are two cases:

- If  $(bc)^+ = (b^+c^+)$ , then  $\omega((bc)^+) = \max(\omega(b^+) + 1, \omega(c^+) + 1)$ , and so we have both:  $\omega(b^+) \leq \omega(b^+) + 1 \leq \max(\omega(b^+) + 1, \omega(c^+) + 1) = \omega((bc)^+)$  and:  $\omega(c^+) \leq \max(\omega(b^+) + 1, \omega(c^+) + 1) = \omega((bc)^+)$ , from which  $(+)$  follows.
- If  $(bc)^+ = (b^+(c^+t))$ , then  $\omega((bc)^+) = \max(\omega(b^+) + 1, \omega(c^+) + 1)$ , and again we have both:  $\omega(b^+) \leq \omega(b^+) + 1 \leq \max(\omega(b^+) + 1, \omega(c^+) + 1) = \omega((bc)^+)$  and:  $\omega(c^+) \leq \omega(c^+) + 1 \leq \max(\omega(b^+) + 1, \omega(c^+) + 1) = \omega((bc)^+)$ , which again implies  $(+)$ . -

Finally, if  $a = s(sb)$ , then:

$$\begin{aligned}
& \omega(a^*) \\
= & \omega((sb)^*(st)) \\
= & \max(\omega(sb)^* + 1, \omega(st)) \\
= & \max(\omega(sb)^* + 1, 1) \\
= & \omega(sb)^* + 1 \\
\leq & \omega(sb)^+ + 1 && \text{by induction} \\
= & \max(\omega(sb)^+ + 1, 1) \\
= & \max(\omega(sb)^+ + 1, \omega(st)) \\
= & \omega((sb)^+(st)) \\
= & \omega((s(sb))^+) \\
= & \omega(a^+)
\end{aligned}$$

... which completes the proof of  $(\Omega)$ .

## A.5 $^\circ$ -Lemma

To be shown:  $(sa)$  is embeddable into  $(sa)^\circ$ , for any  $a \in \mathbf{FT}$ .

We first define a family of closed terms  $\varepsilon_a \in \mathbf{IL}_{(sa)(sa^\circ)}^*$  (where  $a \in \mathbf{FT}$ ) and then show that each  $\varepsilon_a$  is injective, which implies embeddability, given observation  $\mathbf{s}$ ). The following recursion defines the  $\varepsilon$ -terms:

- (a) If  $a = t$  and hence  $(sa)^\circ = st$ ,  $\varepsilon_a := (\lambda p^{\text{st}}. p) [= \kappa_{\text{st}}]$ ;
- (b) if  $a = e$  and hence  $(sa)^\circ = e(st)$ ,  
 $\varepsilon_a := (\lambda f^{\text{se}}. (\lambda x^e. (\lambda i. (f(i) = x)))) [= \kappa_{\text{se}}]$ ;
- (c) if  $a = (sb)$  and hence  $(sa)^\circ = (sb)^\circ(st)$ , two cases must be distinguished:

- (c<sub>1</sub>) if  $b = t$ , then:  $(sb)^o = (st)$  and:  
 $\varepsilon_a := (\lambda C^{s(st)}. (\lambda p^{st}. (\lambda i. (C(i) = p))))$ ;
- (c<sub>2</sub>) if  $b \neq t$ , then:  $(sb)^o = (b^o(st))$  and:  
 $\varepsilon_a := (\lambda C^{s(sb)}. (\lambda y^{(sb)^o}. (\lambda i. (\varepsilon_b(C(i)) = y))))$ ;
- (d) if  $a = (bc)$ , where  $b \neq s$ , and hence  $(sa)^o = (bc)^o(st) = (b^o c^o)(st)$ , four cases must be distinguished:
- (d<sub>1</sub>) if  $b = t = c$ , then:  $(sb)^o = (sc)^o = st$ ,  $(s(bc))^o = (tt)(st)$ , and:  
 $\varepsilon_a := (\lambda f^{s(tt)}. (\lambda h^{tt}. (\lambda i. (f(i) = h))))$ ;
- (d<sub>2</sub>) if  $b = t \neq c$ , then:  $(sb)^o = st$ ,  $(sc)^o = c^o(st)$ ,  $(s(bc))^o = (tc^o)(st)$ ,  
and:  
 $\varepsilon_a := (\lambda f^{s(tc)}. (\lambda h^{tc^o}. (\lambda i. (\forall p^{st}(\forall D^{sc} [(f(i)(p(i)) = D(i)) \leftrightarrow \varepsilon_c(D)(h(p(i))(i))]))$ ;
- (d<sub>3</sub>) if  $b \neq t = c$ , then:  $(sb)^o = b^o(st)$ ,  $(sc)^o = st$ ,  $(s(bc))^o = (b^o t)(st)$ ,  
and:  
 $\varepsilon_a := (\lambda P^{s(bt)}. (\lambda Q^{b^o t}. (\lambda i. (\forall C^{sb})[P(i)(C(i)) \leftrightarrow (\exists x^{b^o o})[\varepsilon_b(C)(x)(i) \wedge Q(x)]]))$
- (d<sub>4</sub>) if  $b \neq t \neq c$ , then  $(sb)^o = b^o(st)$ ,  $(sc)^o = c^o(st)$ ,  $(s(bc))^o = (b^o c^o)(st)$ ,  
and:  
 $\varepsilon_a := (\lambda f^{s(bc)}. (\lambda h^{b^o c^o}. (\lambda i. (\forall C^{sb})(\forall D^{sc} [(f(i)(C(i)) = D(i)) \leftrightarrow (\exists x^{b^o o})[\varepsilon_b(C)(x)(i) \wedge \varepsilon_c(D)(h(x))(i)]]))$

Now, if  $a \in \mathbf{FT} \setminus \{t\}$ ,  $(sa)^o = (a^o(st))$ , and so  $\dot{\varepsilon}_a$  is a ternary relation between objects of type  $a$  ( $\neq t$ ), their (not necessarily unique) Russellian counterparts, and possible worlds. It can be shown that  $(C)$  which counterparts a Fregean intension has for a given world, only depends on its value for that world; that  $(D)$  two Fregean intensions share a Russellian counterpart in a given world only if their values for that world coincide; and that  $(E)$  at any world, every Fregean intension has at least one such  $\dot{\varepsilon}$ -counterpart. In other words, the following facts hold for any Fregean type  $a \neq t$ :

- (C) For any  $\mathbf{X}_1, \mathbf{X}_2 \in D_{sa}$ , and  $\mathbf{w} \in D_s$ :  
if  $\mathbf{X}_1(\mathbf{w}) = \mathbf{X}_2(\mathbf{w})$ , then:  
 $\{z \in D_{a^o} \mid \dot{\varepsilon}_a(\mathbf{X}_1)(z)(\mathbf{w}) = 1\} = \{z \in D_{a^o} \mid \dot{\varepsilon}_a(\mathbf{X}_2)(z)(\mathbf{w}) = 1\}$ .
- (D) For any  $\mathbf{X}_1, \mathbf{X}_2 \in D_{sa}$ , and  $\mathbf{w} \in D_s$ :  
if  $\mathbf{X}_1(\mathbf{w}) \neq \mathbf{X}_2(\mathbf{w})$ , then  
 $\{z \in D_{a^o} \mid \dot{\varepsilon}_a(\mathbf{X}_1)(z)(\mathbf{w}) = 1\} \cap \{z \in D_{a^o} \mid \dot{\varepsilon}_a(\mathbf{X}_2)(z)(\mathbf{w}) = 1\} = \emptyset$ .
- (E) If  $\mathbf{X} \in D_{sa}$  and  $\mathbf{w} \in D_s$ , there is a  $z \in D_{a^o}$  such that  $\dot{\varepsilon}_a(\mathbf{X})(z)(\mathbf{w}) = 1$ .

(D) and (E) together imply that  $\dot{\varepsilon}_a$  is injective. For if  $\mathbf{X}_1, \mathbf{X}_2 \in D_{sa}$  and  $\mathbf{X}_1 \neq \mathbf{X}_2$ , then  $\mathbf{X}_1(\mathbf{w}_0) \neq \mathbf{X}_2(\mathbf{w}_0)$ , for some  $\mathbf{w}_0 \in D_s$ . Hence, putting  $A := \{z \in D_{a^o} \mid \dot{\varepsilon}_a(\mathbf{X}_1)(z)(\mathbf{w}_0) = 1\}$  and  $B := \{z \in D_{a^o} \mid \dot{\varepsilon}_a(\mathbf{X}_2)(z)(\mathbf{w}_0) = 1\}$ , we have:  $A \cap B = \emptyset$ , by (D). But then by (E),  $A \neq \emptyset$ , which means that  $A \neq B$  and hence there must be some  $z_0 \in D_{a^o}$  such that  $\dot{\varepsilon}_a(\mathbf{X}_1)(\mathbf{w}_0)(z_0) \neq$

$\dot{\varepsilon}_a(\mathbf{X}_2)(\mathbf{w}_0)(\mathbf{z}_0)$ , and thus  $\dot{\varepsilon}_a(\mathbf{X}_1) \neq \dot{\varepsilon}_a(\mathbf{X}_2)$ . So for each  $a \in \mathbf{FT} \setminus \{t\}$ ,  $(D)$  and  $(E)$  give rise to  $(Inj)$  – a fact that will come in handy in the proofs below:

$(Inj)$  For any  $\mathbf{X}, \mathbf{Y} \in D_{sa}$ , if  $\dot{\varepsilon}_a(\mathbf{X}) = \dot{\varepsilon}_a(\mathbf{Y})$ , then  $\mathbf{X} = \mathbf{Y}$ .

Since  $\dot{\varepsilon}_t$  is trivially injective,  $(D)$ , and  $(E)$  suffice to establish embeddability of any  $(sa) \in \mathbf{FT}$  into  $(sa)^o \in \mathbf{RT}$ .  $(C)$ ,  $(D)$ , and  $(E)$  can be proved by simultaneous recursion, dropping case  $(a)$ :

$ad(C)$ : Given  $\mathbf{X}_1, \mathbf{X}_2 \in D_{sa}$  and  $\mathbf{w} \in D_s$  such that  $\mathbf{X}_1(\mathbf{w}) = \mathbf{X}_2(\mathbf{w})$ , we show that: for any  $\mathbf{z} \in D_{a^o}$ ,  $\dot{\varepsilon}_a(\mathbf{X}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  $\dot{\varepsilon}_a(\mathbf{X}_2)(\mathbf{z})(\mathbf{w}) = 1$ :

(b) If  $\mathbf{f}_1(\mathbf{w}) = \mathbf{f}_2(\mathbf{w})$ , then by the definition of  $\dot{\varepsilon}_e$ :  
 $\dot{\varepsilon}_e(\mathbf{f}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  $\mathbf{z} = \mathbf{f}_1(\mathbf{w})$  iff  $\mathbf{z} = \mathbf{f}_2(\mathbf{w})$  iff  $\dot{\varepsilon}_e(\mathbf{f}_2)(\mathbf{w}) = 1$ .

(c<sub>1</sub>) The same argument applies in this case.

(c<sub>2</sub>) If  $\mathbf{C}_1(\mathbf{w}) = \mathbf{C}_2(\mathbf{w})$ , we observe:

$\dot{\varepsilon}_{sb}(\mathbf{C}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  
 $\mathbf{z} = \dot{\varepsilon}_b(\mathbf{C}_1(\mathbf{w}))$  iff def.  $\dot{\varepsilon}_{sb}$   
 $\mathbf{z} = \dot{\varepsilon}_b(\mathbf{C}_2(\mathbf{w}))$  iff assumption  
 $\dot{\varepsilon}_{sb}(\mathbf{C}_2)(\mathbf{z})(\mathbf{w}) = 1$ . def.  $\dot{\varepsilon}_{sb}$

(d<sub>1</sub>) The same reasoning as in cases (b) and (c<sub>1</sub>) applies.

(d<sub>2</sub>) Suppose  $\mathbf{f}_1(\mathbf{w}) = \mathbf{f}_2(\mathbf{w})$ . Then for any  $\mathbf{z} \in D_{tc^o}$  we have, by the definition of  $\dot{\varepsilon}_{tc}$ :

$\dot{\varepsilon}_{tc}(\mathbf{f}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  
for all  $\mathbf{p} \in D_{st}$  and  $\mathbf{D} \in D_{sc}$  the following holds:  
 $\mathbf{f}_1(\mathbf{w})(\mathbf{p}(\mathbf{w})) = \mathbf{D}(\mathbf{w})$  iff  $\dot{\varepsilon}_c(\mathbf{D})(\mathbf{z}(\mathbf{p}(\mathbf{w}))) (\mathbf{w}) = 1$ ,  
which means that, for all  $\mathbf{p} \in D_{st}$  and  $\mathbf{D} \in D_{sc}$ :  
 $\mathbf{f}_2(\mathbf{w})(\mathbf{p}(\mathbf{w})) = \mathbf{D}(\mathbf{w})$  iff  $\dot{\varepsilon}_c(\mathbf{D})(\mathbf{z}(\mathbf{p}(\mathbf{w}))) (\mathbf{w}) = 1$ ,  
given that  $\mathbf{f}_1(\mathbf{w}) = \mathbf{f}_2(\mathbf{w})$ . But then, by the definition of  $\dot{\varepsilon}_{tc}$ :  
 $\dot{\varepsilon}_{tc}(\mathbf{f}_2)(\mathbf{z})(\mathbf{w}) = 1$ .

(d<sub>3</sub>) Suppose  $\mathbf{P}_1(\mathbf{w}) = \mathbf{P}_2(\mathbf{w})$ . Then for any  $\mathbf{z} \in D_{b^o t}$  we have:

$\dot{\varepsilon}_{bt}(\mathbf{P}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  
for all  $\mathbf{C} \in D_{sb}$  the following holds:  
 $\mathbf{P}_1(\mathbf{w})(\mathbf{C}(\mathbf{w})) = 1$  iff for some  $\mathbf{x} \in D_{b^o}$ ,  $\dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = \mathbf{z}(\mathbf{x}) = 1$ .  
Hence:  $\dot{\varepsilon}_{bt}(\mathbf{P}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  
for all  $\mathbf{C} \in D_{sb}$  the following holds:  
 $\mathbf{P}_2(\mathbf{w})(\mathbf{C}(\mathbf{w})) = 1$  iff for some  $\mathbf{x} \in D_{b^o}$ ,  $\dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = \mathbf{z}(\mathbf{x}) = 1$ ,  
which means that  $\dot{\varepsilon}_{bt}(\mathbf{P}_1)(\mathbf{z})(\mathbf{w}) = 1$  iff  $\dot{\varepsilon}_{bt}(\mathbf{P}_2)(\mathbf{z})(\mathbf{w}) = 1$ .

(d<sub>4</sub>) Suppose that  $f_1(w) = f_2(w)$ . Then for any  $z \in D_{b^\circ c^\circ}$ :

$$\dot{\varepsilon}_{bc}(f_1)(z)(w) = 1 \text{ iff}$$

for all  $C \in D_{sb}$  and  $D \in D_{sc}$  the following holds:

$$f_1(w)(C(w)) = 1 \text{ iff for some } x \in D_{b^\circ}, \dot{\varepsilon}_b(C)(x)(w) = \dot{\varepsilon}_c(D)(z(x))(w) = 1.$$

$$\text{Hence: } \dot{\varepsilon}_{bc}(f_1)(z)(w) = 1 \text{ iff}$$

for all  $C \in D_{sb}$  and  $D \in D_{sc}$  the following holds:

$$f_2(w)(C(w)) = 1 \text{ iff for some } x \in D_{b^\circ}, \dot{\varepsilon}_b(C)(x)(w) = \dot{\varepsilon}_c(D)(z(x))(w) = 1,$$

which means that  $\dot{\varepsilon}_{bc}(f_1)(z)(w) = 1$  iff  $\dot{\varepsilon}_{bc}(f_2)(z)(w) = 1$ .

$ad(D)$ : Given some  $z \in D_{a^\circ}$  such that:

$$(+) \quad \dot{\varepsilon}_a(X_1)(z)(w) = \dot{\varepsilon}_a(X_2)(z)(w) = 1,$$

we show that  $X_1(w) = X_2(w)$ :

(b) By the definition of  $\dot{\varepsilon}_e$ , (+) implies that  $z = X_1(w) = X_2(w)$ .

(c<sub>1</sub>) The same argument applies as in the previous case.

(c<sub>2</sub>) By the definition of  $\dot{\varepsilon}_{sb}$ , (+) implies that  $z = \dot{\varepsilon}_b(X_1(w)) = \dot{\varepsilon}_b(X_2(w))$ .  
But then  $X_1(w) = X_2(w)$ , because  $b$  satisfies  $(Inj)$ , by induction.

(d<sub>1</sub>) The argument from case (b) applies again.

(d<sub>2</sub>) By the definition of  $\dot{\varepsilon}_{tc}$ , (+) implies that for any  $p \in D_{st}$  and  $D \in D_{sc}$ :  
 $X_1(w)(p(w)) = D(w)$  iff  $\dot{\varepsilon}_c(D)(z(p(w)))(w) = 1$  and:

$$X_2(w)(p(w)) = D(w) \text{ iff } \dot{\varepsilon}_c(D)(z(p(w)))(w) = 1.$$

So, for any such  $p$  and  $D$  we have:

$$X_1(w)(p(w)) = D(w) \text{ iff } X_2(w)(p(w)) = D(w).$$

In particular, for any  $n \in D_t$  and  $y \in D_c$ , we may consider the constant functions  $p_n := \{(w', n) | w' \in D_s\} \in D_{st}$  and  $D_y := \{(w', y) | w' \in D_s\} \in D_{sc}$  to conclude:

$$X_1(w)(p_n(w)) = D_y(w) \text{ iff } X_2(w)(p_n(w)) = D_y(w), \text{ i.e.:}$$

$$X_1(w)(n) = y \text{ iff } X_2(w)(n) = y, \text{ for any } n \in D_t \text{ and } y \in D_c, \text{ which means that } X_1(w) = X_2(w), \text{ by extensionality.}$$

(d<sub>3</sub>) By the definition of  $\dot{\varepsilon}_{bt}$ , (+) implies that for any  $C \in D_{sb}$ :

$$X_1(w)(C(w)) = 1 \text{ iff}$$

there is an  $x_1 \in D_{b^\circ}$  such that:  $\dot{\varepsilon}_b(C)(x_1)(w) = z(x_1) = 1$  and:

$$X_2(w)(C(w)) = 1 \text{ iff}$$

there is an  $x_2 \in D_{b^\circ}$  such that:  $\dot{\varepsilon}_b(C)(x_2)(w) = z(x_2) = 1$ .

So, for any such  $D$  we have:

$$X_1(w)(C(w)) = 1 \text{ iff } X_2(w)(C(w)) = 1.$$

As in the previous case, we may now consider the constant functions  $D_y$  (where  $y \in D_b$ ) to conclude:

$$X_1(w)(C_y(w)) = 1 \text{ iff } X_2(w)(C_y(w)) = 1, \text{ i.e.:}$$

$X_1(w)(y) = 1$  iff  $X_2(w)(y) = 1$ , for any  $y \in D_b$ , and thus:  
 $X_1(w) = X_2(w)$ , by extensionality.

- (d<sub>4</sub>) By the definition of  $\dot{\varepsilon}_{tc}$ , (+) implies that for any  $C \in D_{sb}$  and  $D \in D_{sc}$ :  
 $X_1(w)(C(w)) = D(w)$  iff  
for some  $x_1 \in D_{b^\circ}$ :  $\dot{\varepsilon}_b(C)(x_1)(w) = \dot{\varepsilon}_c(D)(z(x_1))(w) = 1$  and:  
 $X_2(w)(C(w)) = D(w)$  iff  
for some  $x_2 \in D_{b^\circ}$ :  $\dot{\varepsilon}_b(C)(x_2)(w) = \dot{\varepsilon}_c(D)(z(x_2))(w) = 1$ .  
So we have, for any such  $C$  and  $D$ :  
 $X_1(w)(C(w)) = D(w)$  iff  $X_2(w)(C(w)) = D(w)$ .  
As before, we now invoke constant functions  $C_x$  and  $D_y$  (where  $x \in D_b$  and  $y \in D_c$ ) to conclude:  
 $X_1(w)(C_x(w)) = D_y(w)$  iff  $X_2(w)(C_x(w)) = D_y(w)$ , i.e.:  
 $X_1(w)(x) = y$  iff  $X_2(w)(x) = y$ , for any  $x \in D_c$  and  $y \in D_d$ ,  
which means that  $X_1(w) = X_2(w)$ , by extensionality.

ad(E):

- (b) Given  $f \in D_{se}$  and  $w \in D_s$ , we put  $z := f(w)$  and observe that:  
 $\dot{\varepsilon}_e(f)(z)(w) = 1$  iff  $z = f(w)$ , which is true by definition.
- (c<sub>1</sub>) Analogously, we put  $z := C(w)$ , for any  $C \in D_{s(st)}$  and  $w \in D_s$ .
- (c<sub>2</sub>) In a similar vein, given  $C \in D_{s(sb)}$  and  $w \in D_s$ ,  $z := \dot{\varepsilon}_b(C(w))$  fits the bill (and uniquely so).
- (d<sub>1</sub>) As in cases (b) and (c<sub>1</sub>), we put:  $z := f(w)$ , whenever  $f \in D_{s(tt)}$  and  $w \in D_s$ .
- (d<sub>2</sub>) Given  $f \in D_{s(tc)}$  and  $w \in D_s$ , we first apply the inductive hypothesis to  $c$  to observe that, for any  $D \in D_{sb}$ :

$$Y_D := \{y \in D_{c^\circ} \mid \dot{\varepsilon}_c(D)(y)(w) = 1\} \neq \emptyset.$$

So we can select a fixed object  $y_D \in D_{c^\circ}$  from each  $Y_D$  and put:

$$z := \{(p(w), y_D) \mid p \in D_{st}, D \in D_{sc}, \text{ and } f(w)(p(w)) = D(w)\}^{\text{xiv}}$$

We first need to show that  $z \in D_{tc^\circ}$ . So suppose  $n \in D_t$ . Then, putting

$$\begin{aligned} p[n] &:= \{(w, n) \mid w \in D_s\} \in D_{st} \text{ and} \\ D[n] &:= \{(w', f(w')(n)) \mid w' \in D_s\} \in D_{sc}, \end{aligned}$$

we observe that  $D[n](w) = f(w)(n) = f(w)(p[n](w))$ , and thus:

$(n, y_{D[n]}) \in z$ . So  $\text{dom}(z) = D_t$ . As to uniqueness, suppose that  $(n, y), (n, y') \in z$ . Then, for some  $p, p' \in D_{st}$  and  $D, D' \in D_{sc}$ :

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<sup>xiv</sup>To be sure, this is short for:

$\{x \mid \text{there are } n \in D_t, y \in D_{c^\circ}, p \in D_{st} \text{ and } D \in D_{sc} \text{ such that:}$   
 $x = (n, y), n = p(w), y = y_D, \text{ and } f(w)(n) = D(w)\}.$

- (i)  $\mathbf{p}(\mathbf{w}) = \mathbf{p}'(\mathbf{w}) = n$ ; and
- (ii)  $\mathbf{D}(\mathbf{w}) = \mathbf{f}(\mathbf{w})(n) = \mathbf{D}'(\mathbf{w})$ ;
- (iii)  $\mathbf{y} = \mathbf{y}_D$  and  $\mathbf{y}' = \mathbf{y}_{D'}$ .

By induction, (D) applies to  $c$  and so, (ii) implies:  $\mathbf{Y}_D = \mathbf{Y}_{D'}$ , which means that  $\mathbf{y}_D = \mathbf{y}_{D'}$ , and so:  $\mathbf{y} = \mathbf{y}_D = \mathbf{y}_{D'} = \mathbf{y}'$ , by (iii).

Now we know that  $\mathbf{z} \in D_{tc^\circ}$ , we can apply  $\dot{\varepsilon}_{tc}$  to conclude:

$\dot{\varepsilon}_{tc}(\mathbf{f})(\mathbf{z})(\mathbf{w}) = 1$  iff for all  $\mathbf{p} \in D_{st}$  and  $\mathbf{D} \in D_{sc}$ :

(!)  $\mathbf{f}(\mathbf{w})(\mathbf{p}(\mathbf{w})) = \mathbf{D}(\mathbf{w})$  iff  $\dot{\varepsilon}_c(\mathbf{D})(\mathbf{z}(\mathbf{p}(\mathbf{w}))) (\mathbf{w}) = 1$ .

And, indeed, given any  $\mathbf{p} \in D_{st}$  and  $\mathbf{D} \in D_{sc}$ , we can show that they satisfy both directions of (!):

“ $\Rightarrow$ ”: Suppose  $\mathbf{f}(\mathbf{w})(\mathbf{p}(\mathbf{w})) = \mathbf{D}(\mathbf{w})$ . Hence  $\mathbf{z}(\mathbf{p}(\mathbf{w})) = \mathbf{y}_D$ , by the definition of  $\mathbf{z}$ . But  $\mathbf{y}_D \in \mathbf{Y}_D$ , and so:  $\dot{\varepsilon}_c(\mathbf{D})(\mathbf{z}(\mathbf{p}(\mathbf{w}))) (\mathbf{w}) = \dot{\varepsilon}_c(\mathbf{D})(\mathbf{y}_D)(\mathbf{w}) = 1$ .

“ $\Leftarrow$ ”: Suppose that  $\dot{\varepsilon}_c(\mathbf{D})(\mathbf{z}(\mathbf{p}(\mathbf{w}))) (\mathbf{w}) = 1$ . Then  $\mathbf{z}(\mathbf{p}(\mathbf{w})) \in \mathbf{Y}_D$ . Now, by the definition of  $\mathbf{z}$ ,  $\mathbf{z}(\mathbf{p}(\mathbf{w})) = \mathbf{y}_{D'}$ , for some  $\mathbf{D}' \in D_{sc}$ , (#)  $\mathbf{f}(\mathbf{w})(\mathbf{p}(\mathbf{w})) = \mathbf{D}'(\mathbf{w})$ . But then  $\mathbf{z}(\mathbf{p}(\mathbf{w})) \in \mathbf{Y}_D \cap \mathbf{Y}_{D'} \neq \emptyset$ , which means that  $\mathbf{D}(\mathbf{w}) = \mathbf{D}'(\mathbf{w})$ , by (D), inductively applied to  $c$ . So by (#), we conclude that  $\mathbf{f}(\mathbf{w})(\mathbf{p}(\mathbf{w})) = \mathbf{D}(\mathbf{w})$

(d<sub>3</sub>) Given  $\mathbf{P} \in D_{s(bt)}$  and  $\mathbf{w} \in D_s$ , we put

$\mathbf{z} := “\{\mathbf{x} \in D_{b^\circ} \mid \text{for some } \mathbf{D} \in D_{sb}: \mathbf{P}(\mathbf{w})(\mathbf{D}(\mathbf{w})) = \dot{\varepsilon}_b(\mathbf{D})(\mathbf{x})(\mathbf{w}) = 1\}”$ .

We then have:  $\dot{\varepsilon}_{bt}(\mathbf{P})(\mathbf{z})(\mathbf{w}) = 1$  iff for all  $\mathbf{C} \in D_{sb}$ :

(!!)  $\mathbf{P}(\mathbf{w})(\mathbf{C}(\mathbf{w})) = 1$  iff for some  $\mathbf{x} \in D_{b^\circ}$ :  $\dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = \mathbf{z}(\mathbf{x}) = 1$ .

And, indeed, given any  $\mathbf{C} \in D_{sb}$ , we can show that it satisfies both directions of (!!):

“ $\Rightarrow$ ”: Suppose  $\mathbf{P}(\mathbf{w})(\mathbf{C}(\mathbf{w})) = 1$ . By induction, there is an  $\mathbf{x} \in D_{b^\circ}$  such that  $\dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = 1$  and thus  $\mathbf{x}$  trivially satisfies:  $\mathbf{z}(\mathbf{x}) = 1$ , due to the definition of  $\mathbf{z}$ .

“ $\Leftarrow$ ”: Suppose  $\mathbf{x} \in D_{b^\circ}$  and  $\dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = \mathbf{z}(\mathbf{x}) = 1$ . Then, by the definition of  $\mathbf{z}$ :  $\dot{\varepsilon}_b(\mathbf{D})(\mathbf{x})(\mathbf{w}) = 1$ , for some  $\mathbf{D} \in D_{sb}$  such that  $\mathbf{P}(\mathbf{w})(\mathbf{D}(\mathbf{w})) = 1$ . Now, by induction, (D) applies to  $b$ , and so  $\mathbf{C}(\mathbf{w}) = \mathbf{D}(\mathbf{w})$ , given that  $\dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = \dot{\varepsilon}_b(\mathbf{D})(\mathbf{x})(\mathbf{w}) = 1$ . But then we also have:  $\mathbf{P}(\mathbf{w})(\mathbf{C}(\mathbf{w})) = 1$ , given that  $\mathbf{P}(\mathbf{w})(\mathbf{D}(\mathbf{w})) = 1$ .

(d<sub>4</sub>) Given  $\mathbf{f} \in D_{s(bc)}$  and  $\mathbf{w} \in D_s$ , we first apply the inductive hypothesis to  $b$  and  $c$  to observe that, for any  $\mathbf{C} \in D_{sb}$  and  $\mathbf{D} \in D_{sb}$ :

$$\mathbf{X}_C := \{\mathbf{x} \in D_{b^\circ} \mid \dot{\varepsilon}_b(\mathbf{C})(\mathbf{x})(\mathbf{w}) = 1\} \neq \emptyset \text{ and:}$$

$$\mathbf{Y}_D := \{\mathbf{y} \in D_{c^\circ} \mid \dot{\varepsilon}_c(\mathbf{D})(\mathbf{y})(\mathbf{w}) = 1\} \neq \emptyset.$$

So we can select a fixed object  $\mathbf{y}_D \in D_{c^\circ}$  from each  $\mathbf{Y}_D$  and put:

$z := \{(x, y) \in D_{b^\circ} \times D_{c^\circ} \mid \text{there are } C \in D_{sb} \text{ and } D \in D_{sc} \text{ such that:}$   

$$x \in X_C, f(w)(C(w)) = D(w), \text{ and } y = y_D\}$$

$$\cup \{(x, y_0) \mid x \in D_{b^\circ} \setminus \bigcup_{C \in D_{sb}} X_C\},$$
where  $y_0 \in D_{c^\circ}$  is some fixed object. We first need to show that  $z \in D_{b^\circ c^\circ}$ . So suppose  $x \in D_{b^\circ}$ . Then either (Case 1)  $x \notin X_C$ , for any  $C \in D_{sb}$ , and thus  $(x, y_0) \in z$ . Or else (Case 2)  $x \in X_C$ , where  $C \in D_{sb}$ ; then, putting  $D := \{(w', f(w')(C(w')) \mid w' \in D_s\} \in D_{sc}$ , we observe that  $f(w)(C(w)) = D(w)$  and thus  $(x, y_D) \in z$ . So  $\text{dom}(z) = D_{b^\circ}$ . As to uniqueness, suppose that  $(x, y), (x, y') \in z$ . Then if Case 1 holds,  $y = y_0 = y'$ . And if (Case 2) holds, there are  $C, C' \in D_{sb}$  and  $D, D' \in D_{sc}$  such that:

- (i)  $x \in X_C$  and  $x \in X_{C'}$ ; and
- (ii)  $f(w)(C(w)) = D(w)$  and  $f(w)(C'(w)) = D'(w)$ ;
- (iii)  $y = y_D$  and  $y' = y_{D'}$ .

By induction,  $(D)$  applies to  $b$  and so, (i) implies:  $C(w) = C'(w)$ . But then, by (ii):  $D(w) = D'(w)$ . So, applying  $(D)$  to  $c$ , we find:  $Y_D = Y_{D'}$ , which means that  $y_D = y_{D'}$ , and so:  $y = y_D = y_{D'} = y'$ , by (iii).

Now we know that  $z \in D_{b^\circ c^\circ}$ , we can apply  $\dot{\varepsilon}_{bc}$  and conclude:

$\dot{\varepsilon}_{bc}(f)(z)(w) = 1$  iff for all  $C \in D_{sb}$  and  $D \in D_{sc}$ :

(!!!)  $f(w)(C(w)) = D(w)$  iff for some  $x \in D_{b^\circ}$ :

$$\dot{\varepsilon}_b(C)(x)(w) = \dot{\varepsilon}_c(D)(z(x))(w) = 1.$$

And, indeed, given any  $C \in D_{sb}$  and  $D \in D_{sc}$ , we can show that they satisfy both directions of (!!!):

“ $\Rightarrow$ ”: Suppose  $f(w)(C(w)) = D(w)$ . By induction there is some  $x \in D_{b^\circ}$  such that  $\dot{\varepsilon}_b(C)(x)(w) = 1$ , i.e.:  $x \in X_C$ , and thus, by the definition of  $z$ :  $z(x) = y_D \in Y_D$ , whence  $\dot{\varepsilon}_c(D)(z(x))(w) = \dot{\varepsilon}_c(D)(y_D)(w) = 1$ .

“ $\Leftarrow$ ”: Suppose that  $\dot{\varepsilon}_b(C)(x)(w) = \dot{\varepsilon}_c(D)(z(x))(w) = 1$ . Then  $x \in X_C$  and thus, for some  $D' \in D_{sc}$ , we have:  $f(w)(C(w)) = D'(w)$  and  $z(x) = y_{D'}$ . But then  $z(x) \in Y_D \cap Y_{D'}$  and, inductively applying  $(D)$  to  $c$ , we may conclude that  $D(w) = D'(w)$  and thus that  $f(w)(C(w)) = D(w)$ .

## A.6 Translation Theorem

To be shown:  $\llbracket \bar{\alpha} \rrbracket^{F^*, g^*} = \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g})$ , for any  $\alpha \in IL_a^*$  (where  $a \in \mathbf{FT}$ ). The proof proceeds by induction on  $\alpha$  where, without loss of generality, we will rely on the paraphrases given in observation **o**) above:

$$\begin{aligned}
- & \text{ If } \alpha = c(i), \text{ where } c \in Con_{sa}, \text{ then:} \\
& \llbracket \bar{\alpha} \rrbracket^{F^*, g^*} \\
= & \llbracket \tilde{c} \rrbracket^{F^*, g^*} && \text{definition of } \bar{\alpha} \\
= & F^*(\tilde{c}) && \text{interpretation of } IL^* \\
= & \dot{\kappa}_{sa}(F(c)) && \text{definition of } F^* \\
= & \dot{\kappa}_{sa}([\hat{w} : F(c)(w)]) && \eta\text{-conversion} \\
= & \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g}) && \text{interpretation of } IL^*
\end{aligned}$$

$$\begin{aligned}
- & \text{ If } \alpha = x \in Var_a, \text{ we first show that:} \\
& (!) \llbracket (\lambda C^{(sa)^*}. [\Box \Psi_a(C)(\tilde{x})]) \rrbracket^{F^*, g^*} = / \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g})/. \\
& \text{Indeed, for any } C \in D_{(sa)^*}, \text{ we have:} \\
& \llbracket (\lambda C^{(sa)^*}. [\Box \Psi_a(C)(\tilde{x})]) \rrbracket^{F^*, g^*}(C) = 1 \\
\text{iff} & \text{ for all } w \in D_s, \llbracket \Psi_a \rrbracket(C)(g^*(\tilde{x}))(w) = 1 && \text{interpretation of } IL^* \\
\text{iff} & \text{ for all } w \in D_s \text{ there is a } c \in D_{sa} \text{ such that:} \\
& C = \dot{\kappa}_{sa}(c) \text{ and } g^*(\tilde{x}) = \dot{\kappa}_a(c(w)) && \Psi\text{-Lemma} \\
\text{iff} & \text{ for all } w \in D_s \text{ there is a } c \in D_{sa} \text{ such that:} \\
& C = \dot{\kappa}_{sa}(c) \text{ and } \dot{\kappa}_a(g(x)) = \dot{\kappa}_a(c(w)) && \text{def. } g^* \\
\text{iff} & \text{ for all } w \in D_s \text{ there is a } c \in D_{sa} \text{ such that:} \\
& \dot{\kappa}_{sa}^{-1}(C) = c \text{ and } g(x) = c(w) && \text{injectivity of } \dot{\kappa}_{sa} \text{ and } \dot{\kappa}_a \\
\text{iff} & \text{ for all } w \in D_s: g(x) = \dot{\kappa}_{sa}^{-1}(C)(w) && \text{identity law} \\
\text{iff} & \text{ for all } w \in D_s: [\hat{w} : g(x)](w) = \dot{\kappa}_{sa}^{-1}(C)(w) && \beta\text{-conversion} \\
\text{iff} & [\hat{w} : g(x)] = \dot{\kappa}_{sa}^{-1}(C) && \text{extensionality} \\
\text{iff} & \dot{\kappa}_{sa}([\hat{w} : g(x)]) = C && \text{def. } \dot{\kappa}_{sa}^{-1} \\
\text{iff} & \dot{\kappa}_{sa}(\llbracket (\lambda i. x) \rrbracket^{F, g}) = C && IL^*\text{-interpretation}
\end{aligned}$$

Given (!), we thus have:

$$\llbracket \bar{\alpha} \rrbracket^{F^*, g^*} = \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g}).$$

But then  $(sa)^*$  is of the form  $(\underline{a} \text{ (st)})$ , by **i**), and so **q**) applies and we may conclude that:

$$\llbracket \bar{\alpha} \rrbracket^{F^*, g^*} = \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g}).$$

$$\begin{aligned}
- & \text{ If } \alpha = (\beta = \gamma), \text{ where } \beta, \gamma \in IL_b^*, \text{ it needs to be shown that:} \\
& \llbracket (\beta = \gamma) \rrbracket^{F^*, g^*} = \llbracket (\lambda i. (\beta = \gamma)) \rrbracket^{F, g}, \\
& \text{since } \dot{\kappa}_{st} = id_{st}. \text{ And indeed, for any } w \in D_s \text{ we have:} \\
& \llbracket (\beta = \gamma) \rrbracket^{F^*, g^*}(w) = 1 \\
\text{iff} & \text{ there is a } Z \in D_{b^*} \text{ such that:} \\
& \llbracket \Psi_b \rrbracket(\llbracket \bar{\beta} \rrbracket^{F^*, g^*})(Z)(w) = 1 \text{ and } \llbracket \Psi_b \rrbracket(\llbracket \bar{\gamma} \rrbracket^{F^*, g^*})(Z)(w) = 1 && IL^*\text{-int.} \\
\text{iff} & \text{ there is a } Z \in D_{b^*} \text{ such that:} \\
& \llbracket \Psi_b \rrbracket(\dot{\kappa}_{sb}(\llbracket (\lambda i. \beta) \rrbracket^{F, g})(Z)(w)) = \llbracket \Psi_b \rrbracket(\dot{\kappa}_{sb}(\llbracket (\lambda i. \gamma) \rrbracket^{F, g})(Z)(w)) = 1 \text{ ind'n} \\
\text{iff} & \text{ there are } c, d \in D_{sb} \text{ such that:} \\
& \dot{\kappa}_{sb}(\llbracket (\lambda i. \beta) \rrbracket^{F, g}) = \dot{\kappa}_{sb}(c), \dot{\kappa}_{sb}(\llbracket (\lambda i. \gamma) \rrbracket^{F, g}) = \dot{\kappa}_{sb}(d), \\
& \text{and } \dot{\kappa}_b(c(w)) = \dot{\kappa}_b(d(w)) && \Psi\text{-Lemma} \\
\text{iff} & \text{ there are } c, d \in D_{sb} \text{ such that:} \\
& \llbracket (\lambda i. \beta) \rrbracket^{F, g} = c, \llbracket (\lambda i. \gamma) \rrbracket^{F, g} = d \text{ and } c(w) = d(w) && \text{injectivity of } \dot{\kappa}_b \text{ and } \dot{\kappa}_{sb}
\end{aligned}$$

iff  $\llbracket (\lambda i. \beta) \rrbracket^{F,g}(\mathbf{w}) = \llbracket (\lambda i. \gamma) \rrbracket^{F,g}(\mathbf{w})$  identity  
 iff  $\llbracket \beta \rrbracket^{F,g[i/\mathbf{w}]} = \llbracket \gamma \rrbracket^{F,g[i/\mathbf{w}]}$   $IL^*$ -interpretation  
 iff  $\llbracket (\beta = \gamma) \rrbracket^{F,g[i/\mathbf{w}]} = 1$   $IL^*$ -interpretation  
 iff  $\llbracket (\lambda i. (\beta = \gamma)) \rrbracket^{F,g}(\mathbf{w}) = 1$   $IL^*$ -interpretation  
 – If  $\alpha = \beta(\gamma)$ , where  $\beta \in IL_{ba}^*$ ,  $\gamma \in IL_b^*$  and  $b \in \mathbf{FT}$ , then  $\bar{\alpha}$  is of the form  $\bigwedge_{(sa)^*}(\delta)$ , where  $\delta \in IL_{(sa)^*t}^*$ . We first show that:  
 (!! )  $\llbracket \delta \rrbracket^{F^*,g^*} = / \dot{\kappa}_{sa}(\llbracket (\lambda i. \beta(\gamma)) \rrbracket^{F,g}) /$ .  
 Indeed, for any  $\mathbf{C} \in D_{(sa)^*}$ , we have:  
 $\llbracket \delta \rrbracket^{F^*,g^*}(\mathbf{C}) = 1$   
 iff for all  $\mathbf{F} \in D_{(ba)^*}$ ,  $\mathbf{X} \in D_{b^*}$  and  $\mathbf{w} \in D_s$  such that  
 $\llbracket \Psi_{ba} \rrbracket(\llbracket \bar{\beta} \rrbracket^{F^*,g^*})(\mathbf{F})(\mathbf{w}) = \llbracket \Psi_b \rrbracket(\llbracket \bar{\gamma} \rrbracket^{F^*,g^*})(\mathbf{X})(\mathbf{w}) = 1$ ,  
 there is a  $\mathbf{Y} \in D_{a^*}$  such that:  
 $\llbracket \Psi_a \rrbracket(\mathbf{C})(\mathbf{Y})(\mathbf{w}) = \llbracket \Phi_{ba} \rrbracket(\mathbf{F})(\mathbf{X})(\mathbf{Y}) = 1$   $IL^*$ -interpretation  
 iff for all  $\mathbf{F} \in D_{(ba)^*}$ ,  $\mathbf{X} \in D_{b^*}$  and  $\mathbf{w} \in D_s$  such that  
 $\llbracket \Psi_{ba} \rrbracket(\dot{\kappa}_{s(ba)}(\llbracket (\lambda i. \beta) \rrbracket^{F,g}))(\mathbf{F})(\mathbf{w}) = \llbracket \Psi_b \rrbracket(\dot{\kappa}_{sb}(\llbracket (\lambda i. \gamma) \rrbracket^{F,g}))(\mathbf{X})(\mathbf{w}) = 1$ ,  
 there is a  $\mathbf{Y} \in D_{a^*}$  such that:  
 $\llbracket \Psi_a \rrbracket(\mathbf{C})(\mathbf{Y})(\mathbf{w}) = \llbracket \Phi_{ba} \rrbracket(\mathbf{F})(\mathbf{X})(\mathbf{Y}) = 1$  induction  
 iff for all  $\mathbf{F} \in D_{(ba)^*}$ ,  $\mathbf{X} \in D_{b^*}$ ,  $\mathbf{w} \in D_s$ ,  $\mathbf{h} \in D_{s(ba)}$  and  $\mathbf{c} \in D_{sb}$   
 such that  $\dot{\kappa}_{s(ba)}(\llbracket (\lambda i. \beta) \rrbracket^{F,g}) = \dot{\kappa}_{s(ba)}(\mathbf{h})$ ,  $\mathbf{F} = \dot{\kappa}_{ba}(\mathbf{h}(\mathbf{w}))$ ,  
 $\dot{\kappa}_{sb}(\llbracket (\lambda i. \gamma) \rrbracket^{F,g}) = \dot{\kappa}_{sb}(\mathbf{c})$ , and  $\mathbf{X} = \dot{\kappa}_b(\mathbf{c}(\mathbf{w}))$ ,  
 there are  $\mathbf{Y} \in D_{a^*}$ ,  $\mathbf{d} \in D_{sa}$ ,  $\mathbf{g} \in D_{ba}$ , and  $\mathbf{z} \in D_b$  such that:  
 $\mathbf{C} = {}_{sa}(\mathbf{d})$ ,  $\mathbf{Y} = \dot{\kappa}_a(\mathbf{d}(\mathbf{w}))$ ,  $\mathbf{F} = \dot{\kappa}_{ba}(\mathbf{g})$ ,  $\mathbf{X} = \dot{\kappa}_b(\mathbf{z})$ , and  $\mathbf{Y} = \dot{\kappa}_a(\mathbf{g}(\mathbf{z}))$   
 $\Psi$ -/ $\Phi$ -Lemmas  
 iff for all  $\mathbf{F} \in D_{(ba)^*}$ ,  $\mathbf{X} \in D_{b^*}$ ,  $\mathbf{w} \in D_s$ ,  $\mathbf{h} \in D_{s(ba)}$  and  $\mathbf{c} \in D_{sb}$  such that:  
 $\llbracket (\lambda i. \beta) \rrbracket^{F,g} = \mathbf{h}$ ,  $\mathbf{F} = \dot{\kappa}_{ba}(\mathbf{h}(\mathbf{w}))$ ,  $\llbracket (\lambda i. \gamma) \rrbracket^{F,g} = \mathbf{c}$ , and  $\mathbf{X} = \dot{\kappa}_b(\mathbf{c}(\mathbf{w}))$   
 there are  $\mathbf{Y} \in D_{a^*}$ ,  $\mathbf{d} \in D_{sa}$ ,  $\mathbf{g} \in D_{ba}$ , and  $\mathbf{z} \in D_b$  such that:  
 $\mathbf{C} = {}_{sa}(\mathbf{d})$ ,  $\mathbf{Y} = \dot{\kappa}_a(\mathbf{d}(\mathbf{w}))$ ,  $\mathbf{F} = \dot{\kappa}_{ba}(\mathbf{g})$ ,  
 $\mathbf{X} = \dot{\kappa}_b(\mathbf{z})$ , and  $\mathbf{Y} = \dot{\kappa}_a(\mathbf{g}(\mathbf{z}))$   $\dot{\kappa}_{sb}$  injective  
 iff for all  $\mathbf{w} \in D_s$  there are  $\mathbf{d} \in D_{sa}$ ,  $\mathbf{g} \in D_{ba}$ , and  $\mathbf{z} \in D_b$  such that:  
 $\mathbf{C} = {}_{sa}(\mathbf{d})$ ,  $\dot{\kappa}_{ba}(\llbracket (\lambda i. \beta) \rrbracket^{F,g}(\mathbf{w})) = \dot{\kappa}_{ba}(\mathbf{g})$ ,  
 $\dot{\kappa}_b(\llbracket (\lambda i. \gamma) \rrbracket^{F,g}(\mathbf{w})) = \dot{\kappa}_b(\mathbf{z})$ , and  $\dot{\kappa}_a(\mathbf{d}(\mathbf{w})) = \dot{\kappa}_a(\mathbf{g}(\mathbf{z}))$  identity  
 iff for all  $\mathbf{w} \in D_s$  there are  $\mathbf{d} \in D_{sa}$ ,  $\mathbf{g} \in D_{ba}$ , and  $\mathbf{z} \in D_b$  such that:  
 $\dot{\kappa}_{sa}^{-1}(\mathbf{C}) = \mathbf{d}$ ,  $\llbracket (\lambda i. \beta) \rrbracket^{F,g}(\mathbf{w}) = \mathbf{g}$ ,  
 $\llbracket (\lambda i. \gamma) \rrbracket^{F,g}(\mathbf{w}) = \mathbf{z}$ , and  $\mathbf{d}(\mathbf{w}) = \mathbf{g}(\mathbf{z})$   $\dot{\kappa}_{sa}, \dot{\kappa}_{ba}, \dot{\kappa}_b, \dot{\kappa}_a$  injective  
 iff for all  $\mathbf{w} \in D_s$ :  $\dot{\kappa}_{sa}^{-1}(\mathbf{C})(\mathbf{w}) = \llbracket (\lambda i. \beta) \rrbracket^{F,g}(\mathbf{w})(\llbracket (\lambda i. \gamma) \rrbracket^{F,g}(\mathbf{w}))$  identity  
 iff for all  $\mathbf{w} \in D_s$ :  $\dot{\kappa}_{sa}^{-1}(\mathbf{C})(\mathbf{w}) = \llbracket \beta \rrbracket^{F,g[i/\mathbf{w}]}(\llbracket \gamma \rrbracket^{F,g[i/\mathbf{w}]})$   $IL^*$ -interpretation  
 iff for all  $\mathbf{w} \in D_s$ :  $\dot{\kappa}_{sa}^{-1}(\mathbf{C})(\mathbf{w}) = \llbracket \beta(\gamma) \rrbracket^{F,g[i/\mathbf{w}]}$   $IL^*$ -interpretation  
 iff for all  $\mathbf{w} \in D_s$ :  $\dot{\kappa}_{sa}^{-1}(\mathbf{C})(\mathbf{w}) = \llbracket (\lambda i. \beta(\gamma)) \rrbracket^{F,g}(\mathbf{w})$   $IL^*$ -interpretation  
 iff  $\dot{\kappa}_{sa}^{-1}(\mathbf{C}) = \llbracket (\lambda i. \beta(\gamma)) \rrbracket^{F,g}$  extensionality  
 iff  $\mathbf{C} = \dot{\kappa}_{sa}(\llbracket (\lambda i. \beta(\gamma)) \rrbracket^{F,g})$  def.  $\dot{\kappa}_{sa}^{-1}$

Given (!! ), we thus have:

$$\llbracket \bar{\alpha} \rrbracket^{F^*,g^*} = \dot{\bigwedge}_{(sa)^*} (/ \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F,g}) /),$$

from which we may – as above – conclude that:

$$\llbracket \bar{\alpha} \rrbracket^{F^*, g^*} = \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g}).$$

– If  $\alpha = (\lambda x. \beta)$ , where  $x \in Var_b$ ,  $\beta \in IL_b^*$ , and  $a = (bc) \in \mathbf{FT}$ , then  $\bar{\alpha}$  is again of the form  $\sqcap_{(sa)^*}(\delta)$ , where  $\delta \in IL_{(sa)^*t}^*$ . As before, we first show that:

$$(!!!) \quad \llbracket \delta \rrbracket^{F^*, g^*} = / \dot{\kappa}_{sa}(\llbracket (\lambda i. (\lambda x. \beta)) \rrbracket^{F, g})/.$$

Indeed, for any  $\mathbf{G} \in D_{(sa)^*}$ , we have:

$$\llbracket \delta \rrbracket^{F^*, g^*}(\mathbf{G}) = 1$$

iff for all  $\mathbf{X} \in D_{b^*}$  such that  $\llbracket \Gamma_b \rrbracket(\mathbf{X}) = 1$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{F} \in D_{a^*}$  and  $\mathbf{Y} \in D_{c^*}$  such that:

$$\llbracket \Psi_a \rrbracket(\mathbf{G})(\mathbf{F})(\mathbf{w}) = \llbracket \Psi_c \rrbracket(\llbracket \bar{\beta} \rrbracket^{F^*, g^*[\bar{x}/x]})(\mathbf{Y})(\mathbf{w}) =$$

$$\llbracket \Phi_a \rrbracket(\mathbf{F})(\mathbf{X})(\mathbf{Y}) = 1$$

$IL^*$ -interpretation

iff for all  $\mathbf{X} \in rge(\dot{\kappa}_b)$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{F} \in D_{a^*}$  and  $\mathbf{Y} \in D_{c^*}$  such that,

for some  $\mathbf{g} \in D_{sa}$ ,  $\mathbf{f} \in D_a$ ,  $\mathbf{c} \in D_{sc}$ , and  $\mathbf{x} \in D_b$ :

$$\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g}), \mathbf{F} = \dot{\kappa}_a(\mathbf{g}(\mathbf{w})),$$

$$\llbracket \bar{\beta} \rrbracket^{F^*, g^*[\bar{x}/x]} = \dot{\kappa}_{sc}(\mathbf{c}), \mathbf{Y} = \dot{\kappa}_c(\mathbf{c}(\mathbf{w})),$$

$$\mathbf{F} = \dot{\kappa}_a(\mathbf{f}), \mathbf{X} = \dot{\kappa}_b(\mathbf{x}), \text{ and } \mathbf{Y} = \dot{\kappa}_c(\mathbf{f}(\mathbf{x}))$$

$\Psi$ - and  $\Phi$ -Lemmas

iff for all  $\mathbf{X} \in rge(\dot{\kappa}_b)$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{F} \in D_{a^*}$  and  $\mathbf{Y} \in D_{c^*}$  such that:

for some  $\mathbf{g} \in D_{sa}$ ,  $\mathbf{f} \in D_a$ ,  $\mathbf{c} \in D_{sb}$ , and  $\mathbf{x} \in D_b$ :

$$\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g}), \mathbf{F} = \dot{\kappa}_a(\mathbf{f}), \mathbf{g}(\mathbf{w}) = \mathbf{f},$$

$$\llbracket \bar{\beta} \rrbracket^{F^*, g^*[\bar{x}/x]} = \dot{\kappa}_{sc}(\mathbf{c}), \mathbf{Y} = \dot{\kappa}_c(\mathbf{y}), \mathbf{c}(\mathbf{w}) = \mathbf{y},$$

$$\mathbf{X} = \dot{\kappa}_b(\mathbf{x}), \text{ and } \mathbf{f}(\mathbf{x}) = \mathbf{y}$$

injectivity of  $\dot{\kappa}_c$  and  $\dot{\kappa}_a$

iff for all  $\mathbf{x} \in D_b$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{g} \in D_{sa}$ ,  $\mathbf{f} \in D_a$  and  $\mathbf{c} \in D_{sb}$  such that:

$$\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g}), \mathbf{g}(\mathbf{w}) = \mathbf{f}, \llbracket \bar{\beta} \rrbracket^{F^*, g^*[\bar{x}/x]} = \dot{\kappa}_{sc}(\mathbf{c}), \text{ and } \mathbf{f}(\mathbf{x}) = \mathbf{c}(\mathbf{w})$$

identity

iff for all  $\mathbf{x} \in D_b$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{g} \in D_{sa}$ ,  $\mathbf{f} \in D_a$  and  $\mathbf{c} \in D_{sb}$  such that:

$$\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g}), \mathbf{g}(\mathbf{w}) = \mathbf{f}, \llbracket \bar{\beta} \rrbracket^{F^*, g^*[\bar{x}/x]} = \dot{\kappa}_{sc}(\mathbf{c}), \text{ and } \mathbf{f}(\mathbf{x}) = \mathbf{c}(\mathbf{w})$$

n)

iff for all  $\mathbf{x} \in D_b$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{g} \in D_{sa}$ ,  $\mathbf{f} \in D_a$  and  $\mathbf{c} \in D_{sb}$  such that:

$$\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g}), \mathbf{g}(\mathbf{w}) = \mathbf{f}, \dot{\kappa}_{sc}(\llbracket (\lambda i. \beta) \rrbracket^{F, g^*[\bar{x}/x]}) = \dot{\kappa}_{sc}(\mathbf{c}), \text{ and } \mathbf{f}(\mathbf{x}) = \mathbf{c}(\mathbf{w})$$

induction

iff for all  $\mathbf{x} \in D_b$  and all  $\mathbf{w} \in D_s$

there are  $\mathbf{g} \in D_{sa}$ ,  $\mathbf{f} \in D_a$  and  $\mathbf{c} \in D_{sb}$  such that:

$$\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g}), \mathbf{g}(\mathbf{w}) = \mathbf{f}, \llbracket (\lambda i. \beta) \rrbracket^{F, g^*[\bar{x}/x]} = \mathbf{c}, \text{ and } \mathbf{f}(\mathbf{x}) = \mathbf{c}(\mathbf{w})$$

injectivity of  $\dot{\kappa}_{sc}$

iff for all  $\mathbf{x} \in D_b$  and all  $\mathbf{w} \in D_s$  there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$

$$\text{and } \mathbf{g}(\mathbf{w})(\mathbf{x}) = \llbracket (\lambda i. \beta) \rrbracket^{F, g^*[\bar{x}/x]}(\mathbf{w})$$

identity

iff for all  $\mathbf{x} \in D_b$  and all  $\mathbf{w} \in D_s$  there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$

$$\text{and } \mathbf{g}(\mathbf{w})(\mathbf{x}) = \llbracket \beta \rrbracket^{F, g^*[\bar{x}/x]}[\bar{i}/w]$$

$IL^*$ -interpretation

iff for all  $x \in D_b$  and all  $w \in D_s$  there is a  $g \in D_{sa}$  such that  $G = \dot{\kappa}_{sa}(g)$   
 and  $g(w)(x) = \llbracket \beta \rrbracket^{F, g[i/w]}[x/x]$  since  $i \neq x$ ,  $g[x/x][i/w] = g[i/w][x/x]$   
 iff for all  $x \in D_b$  and all  $w \in D_s$  there is a  $g \in D_{sa}$  such that  $G = \dot{\kappa}_{sa}(g)$   
 and  $g(w)(x) = \llbracket (\lambda x. \beta) \rrbracket^{F, g[i/w]}(x)$   $IL^*$ -interpretation  
 iff for all  $x \in D_b$  and all  $w \in D_s$  there is a  $g \in D_{sa}$  such that  $\dot{\kappa}_{sa}^{-1}(G) = g$   
 and  $g(w)(x) = \llbracket (\lambda x. \beta) \rrbracket^{F, g[i/w]}(x)$  injectivity of  $\dot{\kappa}_{sa}$   
 iff for all  $x \in D_b$  and all  $w \in D_s$ :  
 $\dot{\kappa}_{sa}^{-1}(G)(w)(x) = \llbracket (\lambda x. \beta) \rrbracket^{F, g[i/w]}(w)(x)$  identity  
 iff for all  $x \in D_b$  and all  $w \in D_s$ :  
 $\dot{\kappa}_{sa}^{-1}(G)(w)(x) = \llbracket (\lambda i. (\lambda x. \beta)) \rrbracket^{F, g}(w)(x)$   $IL^*$ -interpretation  
 iff  $\dot{\kappa}_{sa}^{-1}(G) = \llbracket (\lambda i. (\lambda x. \beta)) \rrbracket^{F, g}$  extensionality  
 iff  $G = \dot{\kappa}_{sa}(\llbracket (\lambda i. (\lambda x. \beta)) \rrbracket^{F, g})$  def.  $\dot{\kappa}_{sa}$   
 ... which establishes (!!!). As before, we may now apply **s**) and conclude that:

$$\llbracket \bar{\alpha} \rrbracket^{F^*, g^*} = \dot{\kappa}_{sa}(\dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g}) / ) = \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F, g}).$$

– If  $\alpha = \beta(i)$ , where  $\beta \in IL_{sa}^*$ ,  $\bar{\alpha}$  is again of the form  $\sqcap_{(sa)^*}(\delta)$ , where  $\delta \in IL_{(sa)^*t}^*$ ; and again we can show that:

$$(!^4) \quad \llbracket \delta \rrbracket^{F^*, g^*} = \dot{\kappa}_{sa}(\llbracket (\lambda i. \beta(\gamma)) \rrbracket^{F, g}) / .$$

For we have, for any  $G \in D_{(sa)^*}$ :

$$\llbracket \delta \rrbracket^{F^*, g^*}(G) = 1$$

iff for all  $w \in D_s$  there are  $H \in D_{(sa)^*}$  and  $X \in D_{a^*}$  such that:

$$\llbracket \Psi_{sa} \rrbracket(\llbracket \bar{\beta} \rrbracket^{F^*, g^*})(H)(w) = \llbracket \Psi_a \rrbracket(H)(X)(w) = \llbracket \Psi_a \rrbracket(G)(X)(w) = 1$$

$IL^*$ -interpretation

iff for all  $w \in D_s$  there are  $H \in D_{(sa)^*}$ ,  $X \in D_{a^*}$ ,

$b \in D_{s(sa)}$ ,  $g, h \in D_{sa}$  such that:

$$\llbracket \bar{\beta} \rrbracket^{F^*, g^*} = \dot{\kappa}_{s(sa)}(b), H = \dot{\kappa}_{sa}(b(w)),$$

$$H = \dot{\kappa}_{sa}(h), X = \dot{\kappa}_a(h(w)),$$

$$G = \dot{\kappa}_{sa}(g), \text{ and } X = \dot{\kappa}_a(g(w))$$

$\Psi$ -Lemma

iff for all  $w \in D_s$  there are  $b \in D_{s(sa)}$ ,  $g, h \in D_{sa}$  such that:

$$\llbracket \bar{\beta} \rrbracket^{F^*, g^*} = \dot{\kappa}_{s(sa)}(b), \dot{\kappa}_{sa}(b(w)) = \dot{\kappa}_{sa}(h), \dot{\kappa}_a(h(w)) = \dot{\kappa}_a(g(w)),$$

$$\text{and } G = \dot{\kappa}_{sa}(g)$$

identity

iff there is a  $g \in D_{sa}$  such that  $G = \dot{\kappa}_{sa}(g)$  and

for all  $w \in D_s$  there are  $b \in D_{s(sa)}$  and  $h \in D_{sa}$  such that:

$$\llbracket \bar{\beta} \rrbracket^{F^*, g^*} = \dot{\kappa}_{s(sa)}(b), b(w) = h, \text{ and } h(w) = g(w)$$

injectivity of  $\dot{\kappa}_{sa}, \dot{\kappa}_a$

iff there is a  $g \in D_{sa}$  such that  $G = \dot{\kappa}_{sa}(g)$  and

for all  $w \in D_s$  there is a  $b \in D_{s(sa)}$  such that:

$$\llbracket \bar{\beta} \rrbracket^{F^*, g^*} = \dot{\kappa}_{s(sa)}(b) \text{ and } b(w)(w) = g(w)$$

identity

iff there is a  $g \in D_{sa}$  such that  $G = \dot{\kappa}_{sa}(g)$  and

for all  $w \in D_s$  there is a  $b \in D_{s(sa)}$  such that:

$$\dot{\kappa}_{s(sa)}(\llbracket (\lambda i. \beta) \rrbracket^{F, g}) = \dot{\kappa}_{s(sa)}(b) \text{ and } b(w)(w) = g(w)$$

induction

iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$  and  
 for all  $\mathbf{w} \in D_s$  there is a  $\mathbf{b} \in D_{s(sa)}$  such that:  
 $\llbracket (\lambda i. \beta) \rrbracket^{F,g} = \mathbf{b}$  and  $\mathbf{b}(\mathbf{w})(\mathbf{w}) = \mathbf{g}(\mathbf{w})$  injectivity of  $\dot{\kappa}_{s(sa)}$   
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$   
 and for all  $\mathbf{w} \in D_s$ :  $\llbracket (\lambda i. \beta) \rrbracket^{F,g}(\mathbf{w})(\mathbf{w}) = \mathbf{g}(\mathbf{w})$  identity  
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$   
 and for all  $\mathbf{w} \in D_s$ :  $\llbracket \beta \rrbracket^{F,g[i/\mathbf{w}]}(\mathbf{w}) = \mathbf{g}(\mathbf{w})$   $IL^*$ -interpretation  
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$   
 and for all  $\mathbf{w} \in D_s$ :  $\llbracket \beta \rrbracket^{F,g[i/\mathbf{w}]}(\mathbf{g}[i/\mathbf{w}](i)) = \mathbf{g}(\mathbf{w})$  def.  $\mathbf{g}[i/\mathbf{w}]$   
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$   
 and for all  $\mathbf{w} \in D_s$ :  $\llbracket \beta \rrbracket^{F,g[i/\mathbf{w}]}(\llbracket i \rrbracket^{F,g[i/\mathbf{w}]}) = \mathbf{g}(\mathbf{w})$   $IL^*$ -interpretation  
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$   
 and for all  $\mathbf{w} \in D_s$ :  $\llbracket \beta(i) \rrbracket^{F,g[i/\mathbf{w}]} = \mathbf{g}(\mathbf{w})$   $IL^*$ -interpretation  
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$   
 and for all  $\mathbf{w} \in D_s$ :  $\llbracket \lambda i. \beta(i) \rrbracket^{F,g}(\mathbf{w}) = \mathbf{g}(\mathbf{w})$   $IL^*$ -interpretation  
 iff there is a  $\mathbf{g} \in D_{sa}$  such that  $\mathbf{G} = \dot{\kappa}_{sa}(\mathbf{g})$  and  $\llbracket \lambda i. \beta(i) \rrbracket^{F,g} = \mathbf{g}$  extensionality  
 iff  $\mathbf{G} = \dot{\kappa}_{sa}(\llbracket \lambda i. \beta(i) \rrbracket^{F,g})$  identity  
 ... and we may again conclude that:  
 $\llbracket \bar{\alpha} \rrbracket^{F^*,g^*} = \llbracket \sqcap_{sa} \rrbracket / \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F,g}) / = \dot{\kappa}_{sa}(\llbracket (\lambda i. \alpha) \rrbracket^{F,g})$ .

– If  $\alpha = (\lambda i. \beta)$ , where  $\beta \in IL_b^*$ , then  $\bar{\alpha}$  is of the form  $\sqcap(\delta)$  again, where  $\delta \in IL_{(sa)^*t}^*$ ; and once more we can show that:

$$(!^5) \quad \llbracket \delta \rrbracket^{F^*,g^*} = / \dot{\kappa}_{sa}(\llbracket (\lambda i. (\lambda i. \beta)) \rrbracket^{F,g}) /.$$

For we have, for any  $\mathbf{C} \in D_{(sa)^*}$ :

$\llbracket (\lambda C^{(sa)^*}. \sqcap \Psi_a(C)(\bar{\beta})) \rrbracket^{M^*,g^*}(\mathbf{C}) = 1$   
 iff for all  $\mathbf{w} \in D_s$ :  $\llbracket \Psi_a \rrbracket(\mathbf{C})(\llbracket \bar{\beta} \rrbracket^{M^*,g^*})(\mathbf{w}) = 1$   $IL^*$ -interpretation  
 iff for all  $\mathbf{w} \in D_s$  there is a  $\mathbf{c} \in D_{sa}$  such that:  
 $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$  and  $\llbracket \bar{\beta} \rrbracket^{M^*,g^*} = \dot{\kappa}_a(\mathbf{c}(\mathbf{w}))$   $\Psi$ -Lemma  
 iff for all  $\mathbf{w} \in D_s$  there is a  $\mathbf{c} \in D_{sa}$  such that:  
 $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$  and  $\dot{\kappa}_a(\llbracket (\lambda i. \beta) \rrbracket^{M,g}) = \dot{\kappa}_a(\mathbf{c}(\mathbf{w}))$  induction  
 iff for all  $\mathbf{w} \in D_s$  there is a  $\mathbf{c} \in D_{sa}$  such that:  
 $\dot{\kappa}_{sa}^{-1}(\mathbf{C}) = \mathbf{c}$  and  $\llbracket (\lambda i. \beta) \rrbracket^{M,g} = \mathbf{c}(\mathbf{w})$  injectivity of  $\dot{\kappa}_{sa}$  and  $\dot{\kappa}_a$   
 iff there is a  $\mathbf{c} \in D_{sa}$  such that  $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$  and  
 for all  $\mathbf{w} \in D_s$ :  $\llbracket (\lambda i. \beta) \rrbracket^{M,g} = \mathbf{c}(\mathbf{w})$  injectivity of  $\dot{\kappa}_a$   
 iff there is a  $\mathbf{c} \in D_{sa}$  such that  $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$  and  
 for all  $\mathbf{w} \in D_s$ :  $\llbracket (\lambda i. \beta) \rrbracket^{M,g[i/\mathbf{w}]} = \mathbf{c}(\mathbf{w})$   $i \notin Fr((\lambda i. \beta))$   
 iff there is a  $\mathbf{c} \in D_{sa}$  such that  $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$  and  
 for all  $\mathbf{w} \in D_s$ :  $\llbracket (\lambda i. (\lambda i. \beta)) \rrbracket^{M,g}(\mathbf{w}) = \mathbf{c}(\mathbf{w})$   $IL^*$ -interpretation  
 iff there is a  $\mathbf{c} \in D_{sa}$  such that  $\mathbf{C} = \dot{\kappa}_{sa}(\mathbf{c})$  and  
 $\llbracket (\lambda i. (\lambda i. \beta)) \rrbracket^{M,g} = \mathbf{c}$  extensionality  
 iff  $\mathbf{C} = \dot{\kappa}_{sa}(\llbracket (\lambda i. (\lambda i. \beta)) \rrbracket^{M,g})$  identity  
 Now  $(!^5)$  follows as before.

## References

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